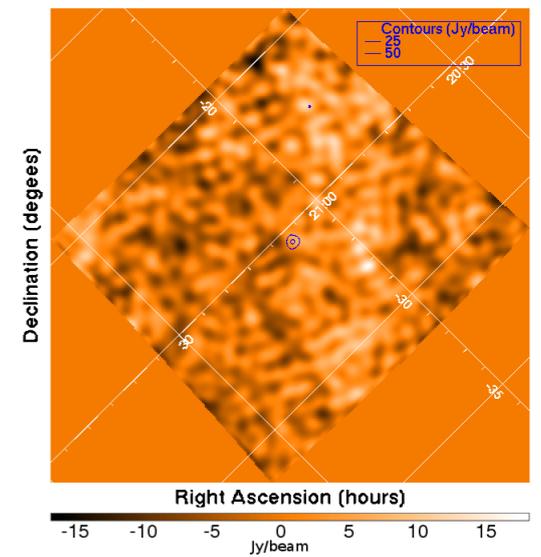
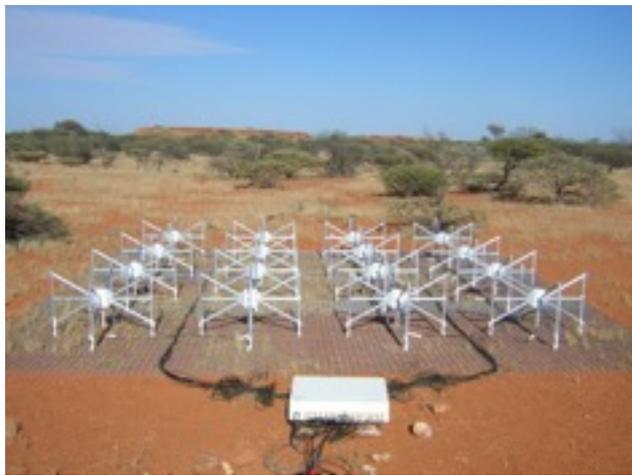


# Using GPUs for Signal Correlation



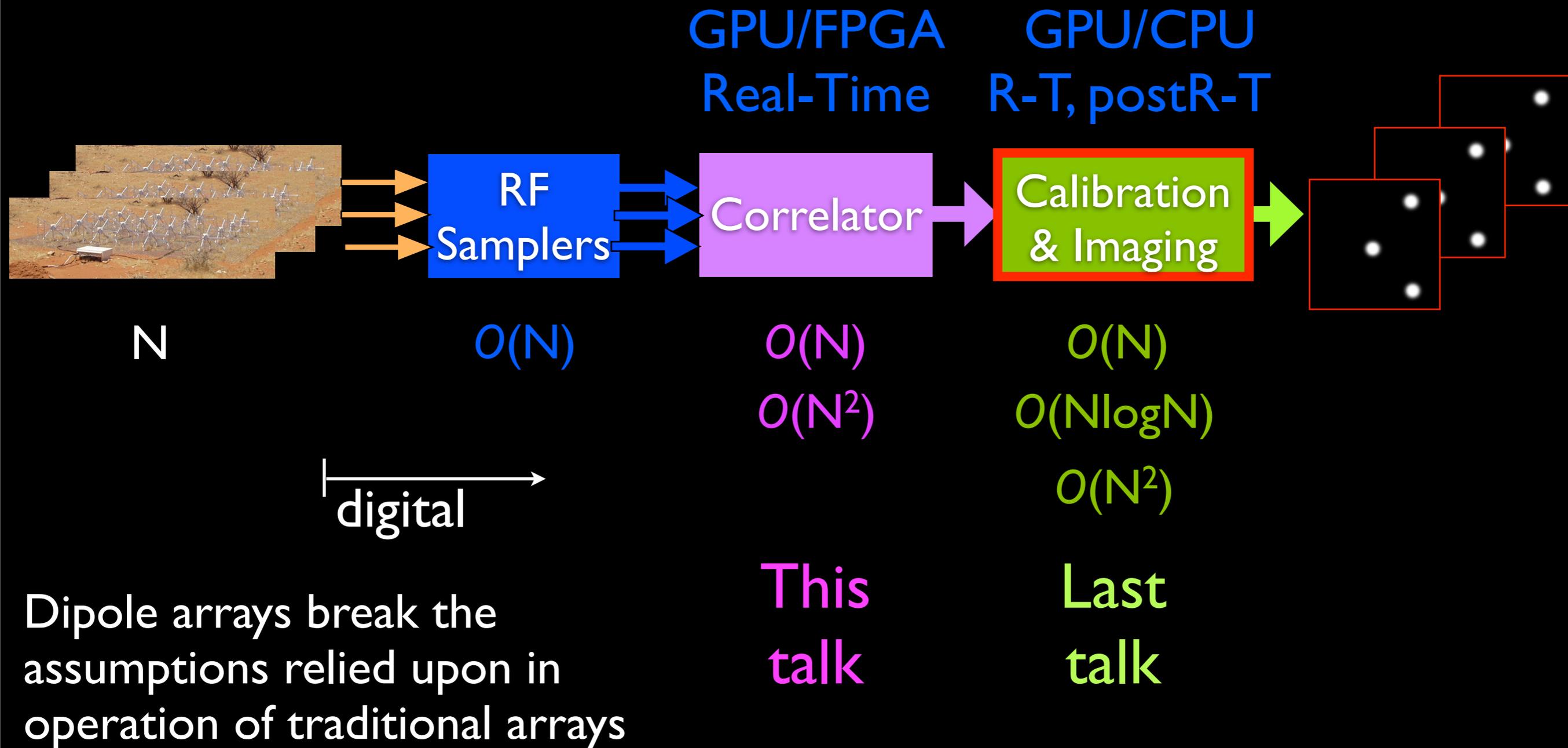
Michael Clark  
with  
Paul La Plante and Lincoln Greenhill

# Outline

- Motivation
- Mapping the *X*-engine onto a GPU
  - Or how to get a sustained TFLOP from a Fermi
- Comparison
- Summary and Conclusions

# Dipole Array Signal Processing

Heterogeneous HPC solution



# Correlator

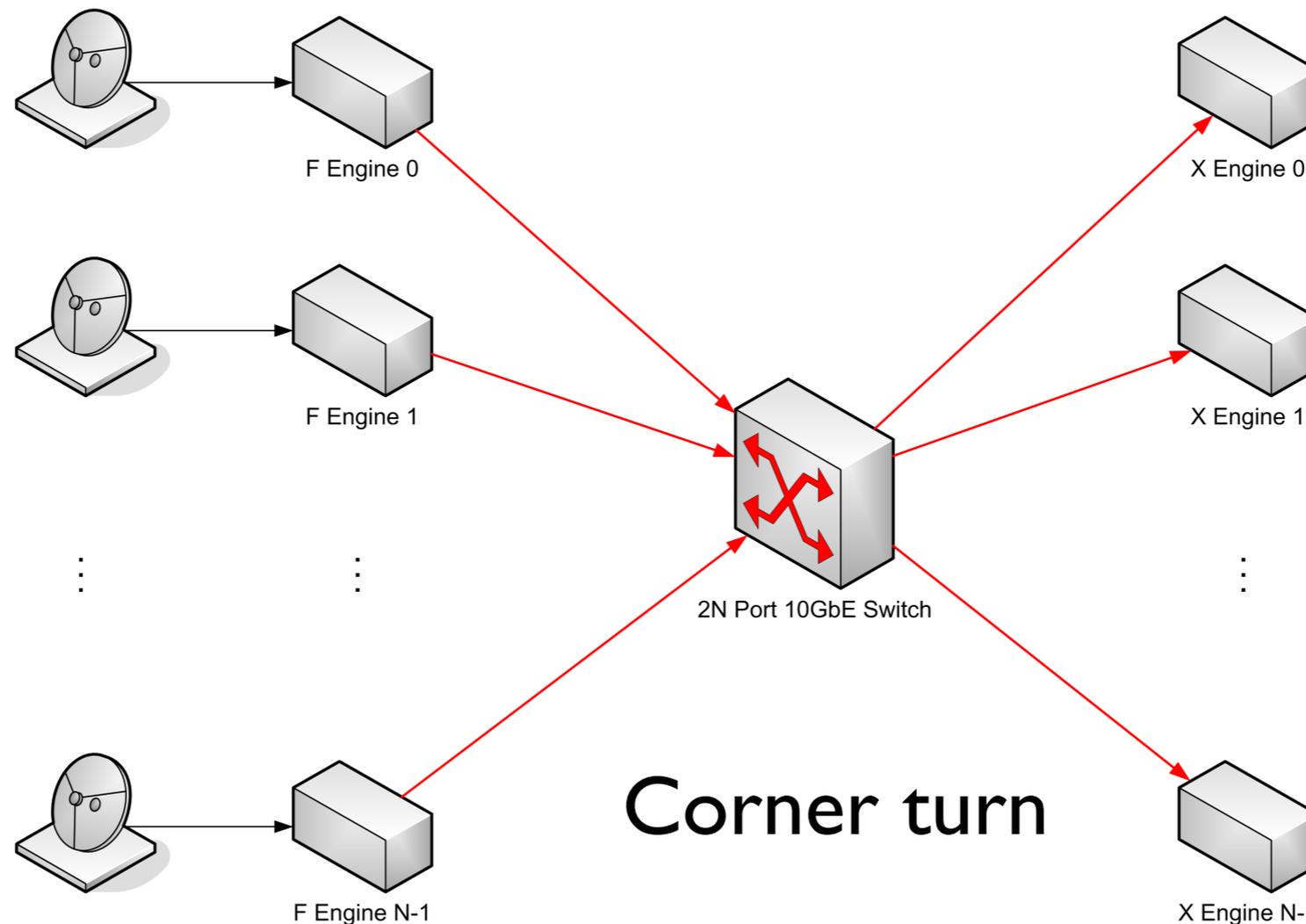
- Builds the cross-power spectrum of the sky

$$S_{ij}(\nu) = \int_{-\infty}^{\infty} (A_i \star A_j)(\tau) e^{-i2\pi\nu\tau} d\tau$$

- Cross-correlate the signal from every station pair
- XF Correlator
  - Cross correlate the signal then FFT
- FX Correlator
  - From convolution theorem  $\mathcal{F}(A \star B) = (\mathcal{F}A) \times (\mathcal{F}B)$
  - FFT the signal then cross-multiply  $O(B N_s (N_s + k))$

# FX Correlator

(Figure taken from Casper collaboration)

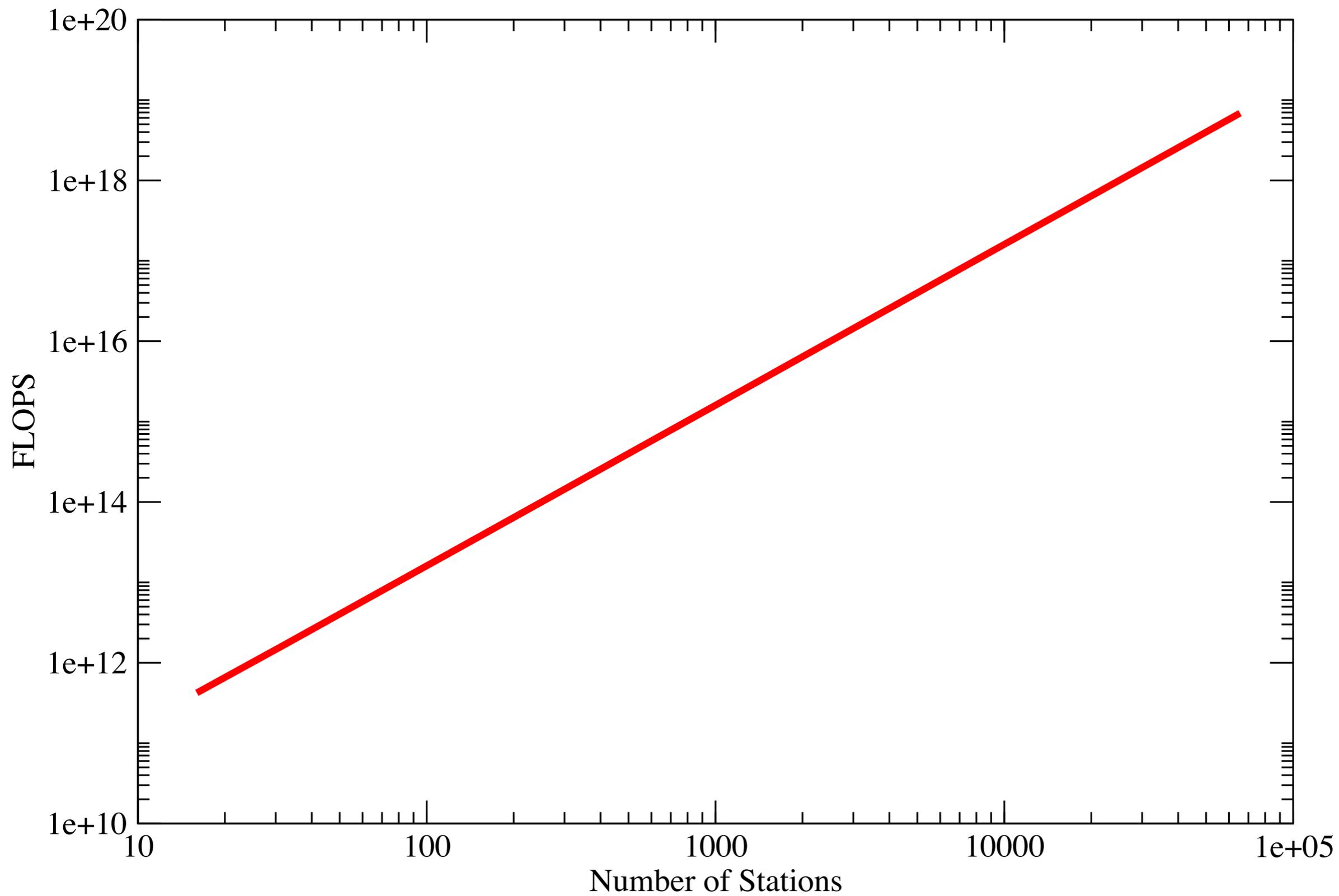


F-engine compute  
scales linearly with  $N_s$

X-engine compute  
scale quadratically with  $N_s$

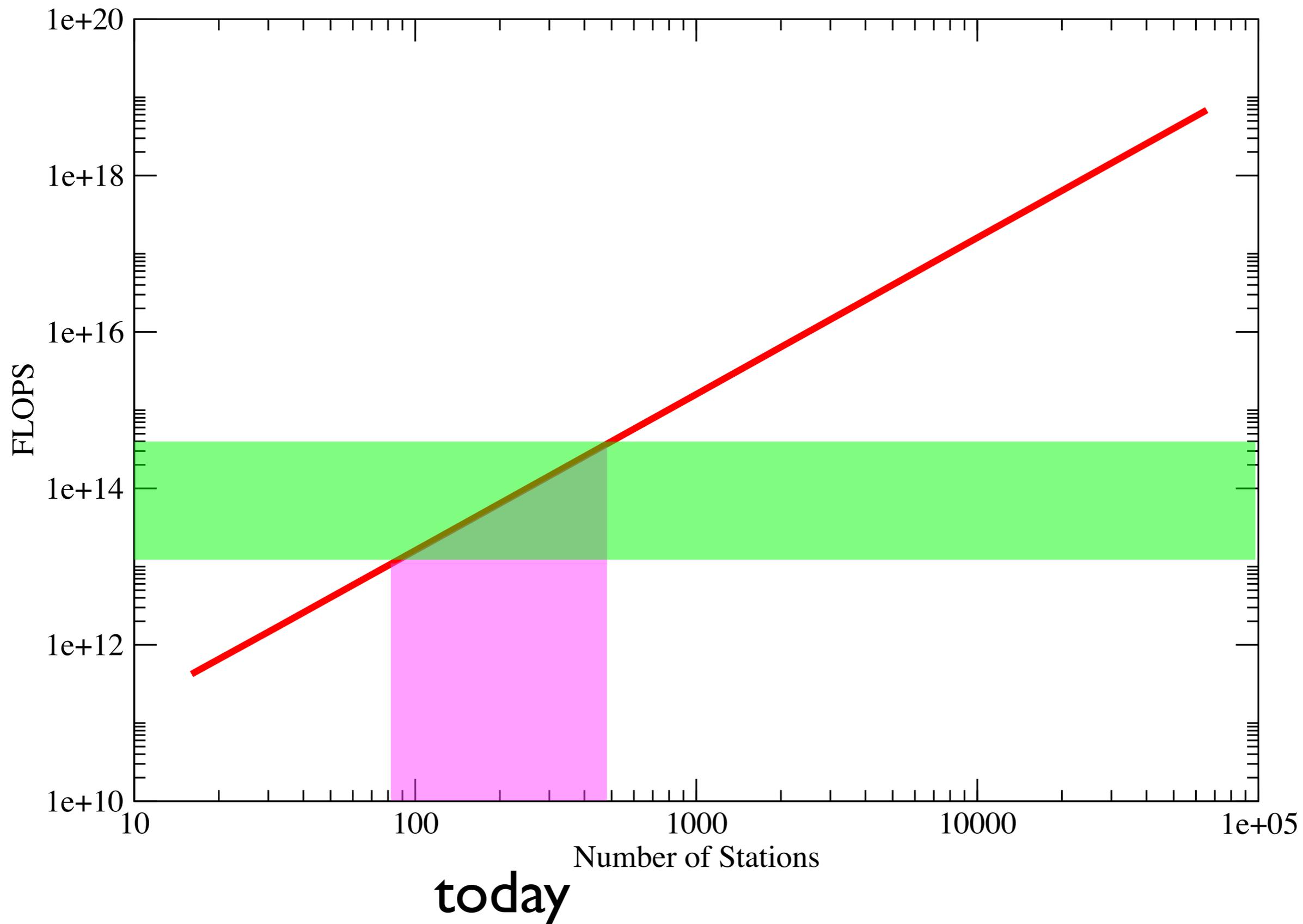
# X-Engine Computing Requirements

$N_c = 1000, B = 100 \text{ MHz}$



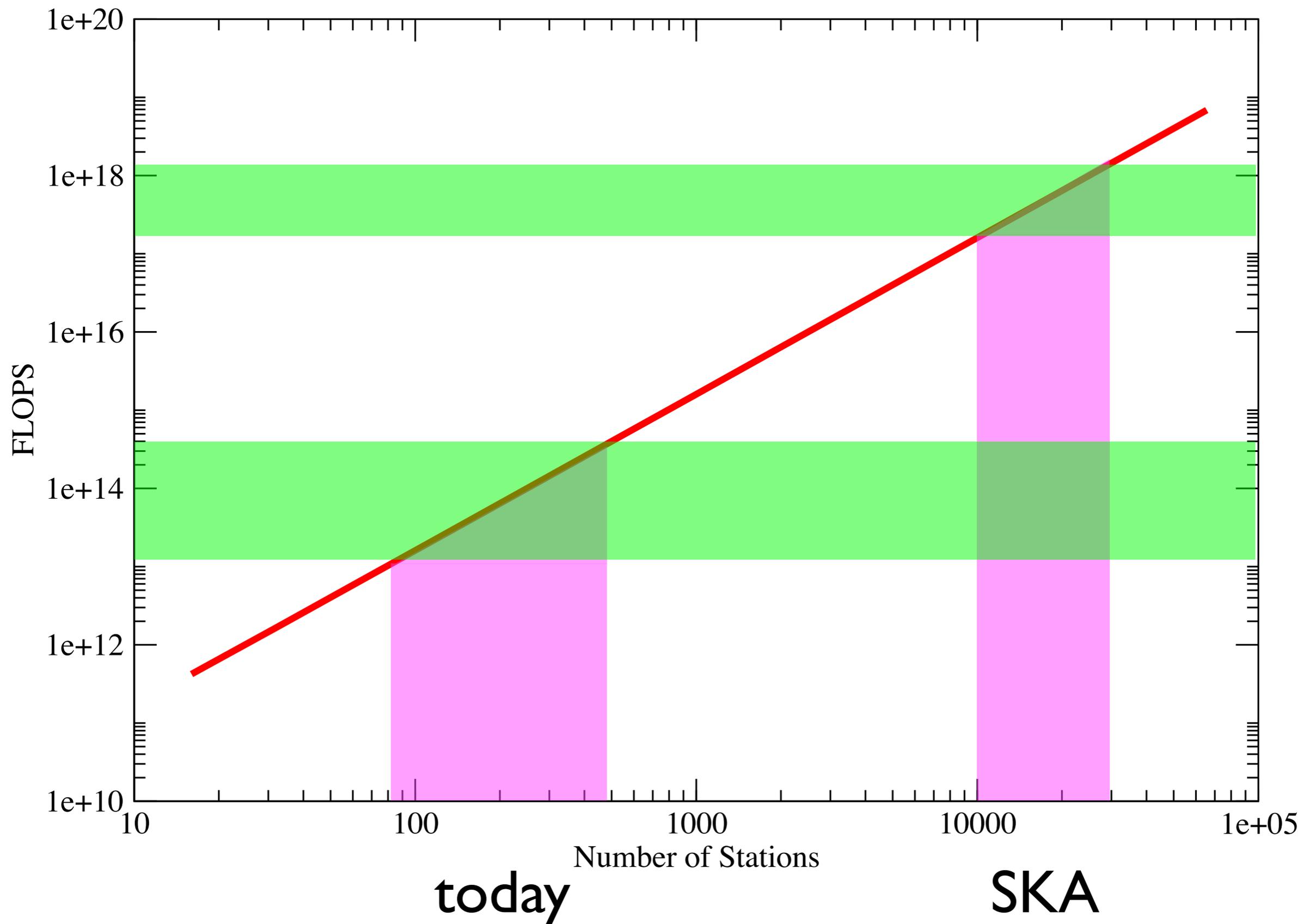
# X-Engine Computing Requirements

$$N_c = 1000, B = 100 \text{ MHz}$$



# X-Engine Computing Requirements

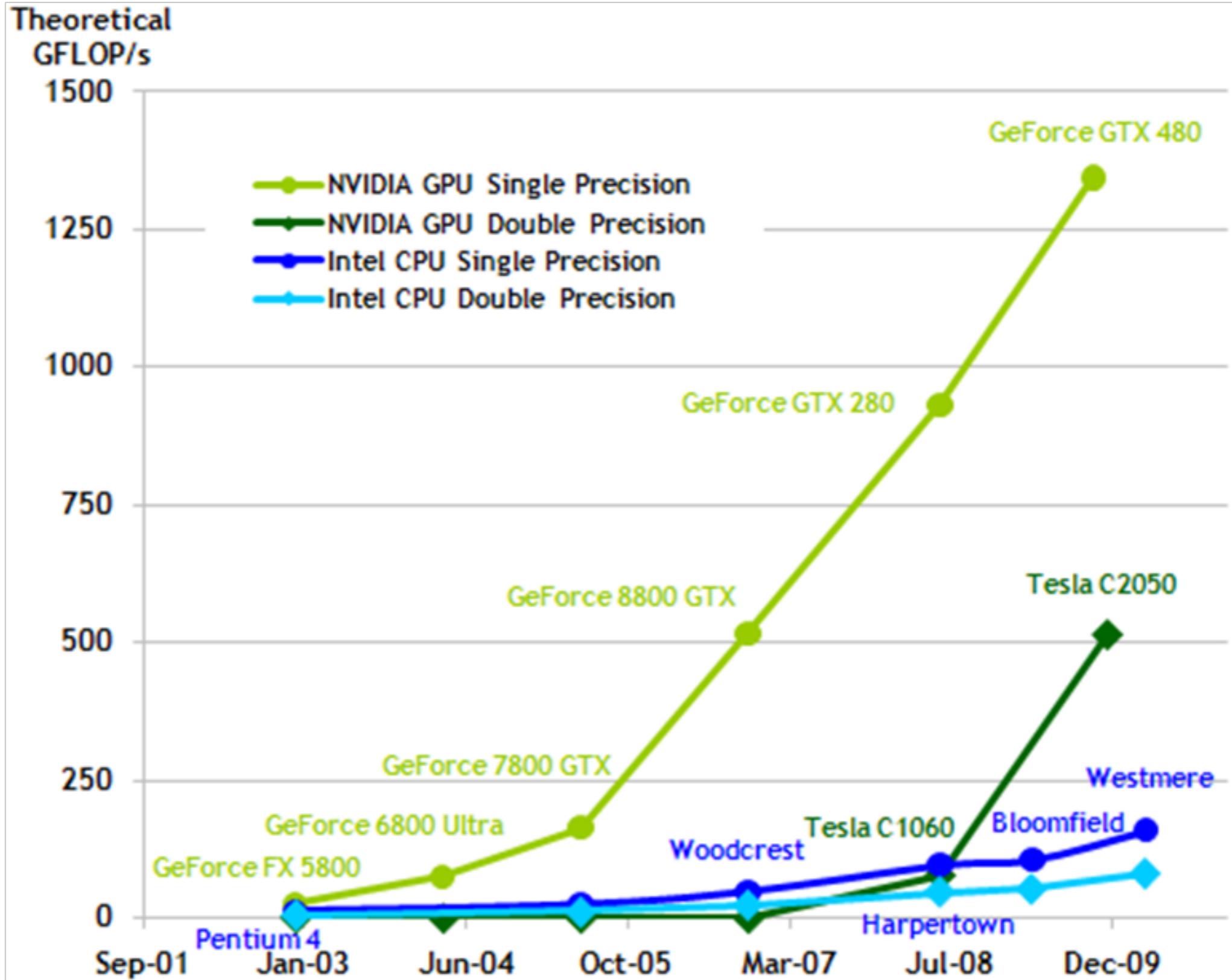
$$N_c = 1000, B = 100 \text{ MHz}$$



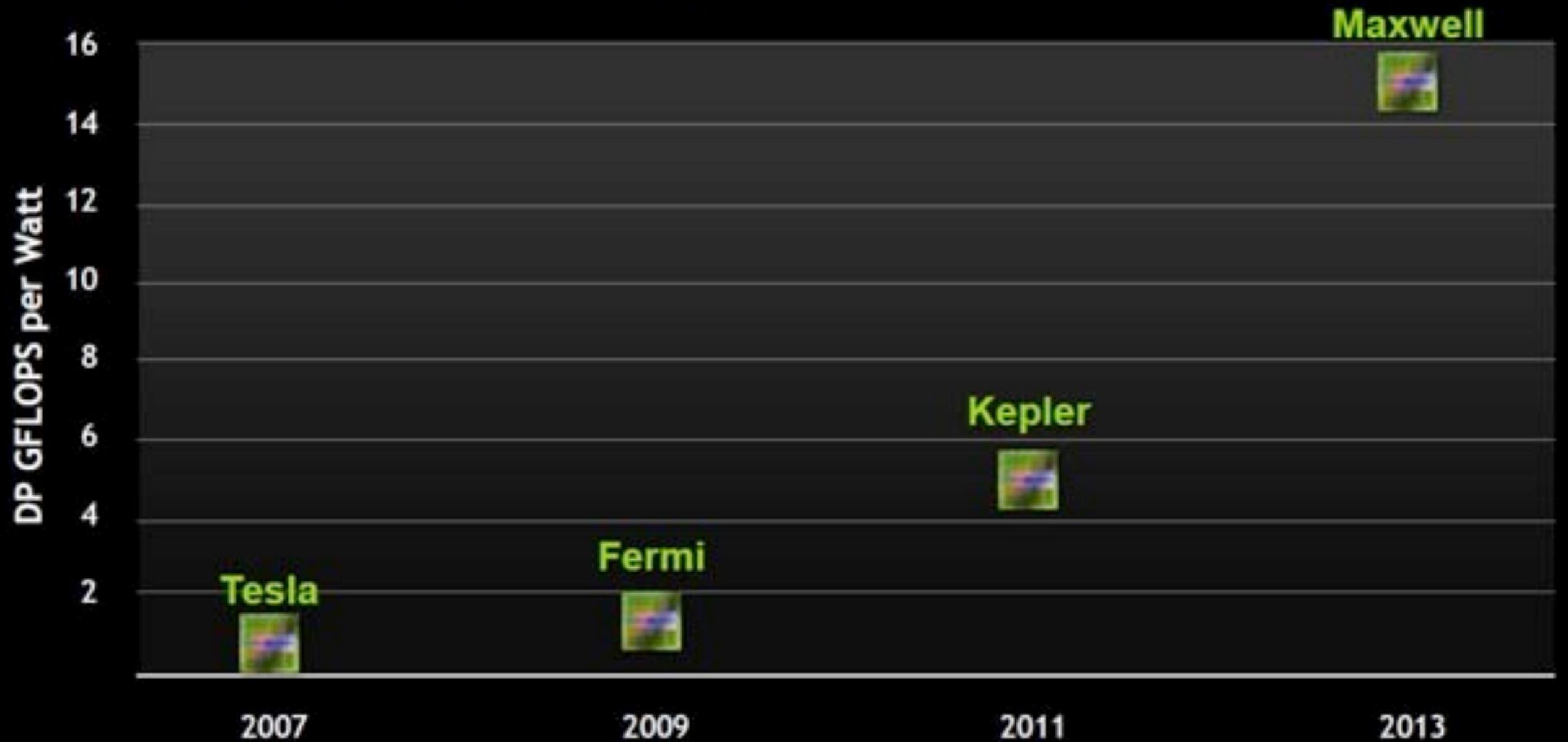
# FPGAs

- FPGAs (or ASICs) generally used for correlation
- Ideal since only require limited fixed-point ops
- Very power efficient since all die area devoted to problem
- Extremely expensive
  - Astronomers survive on donations and old h/w
- Development time can be very long
- Local experts at Berkeley

[http://casper.berkeley.edu/wiki/CASPER\\_Correlator](http://casper.berkeley.edu/wiki/CASPER_Correlator)



# CUDA GPU Roadmap



Huang's GTC Keynote

**Are GPUs (part of) the solution?**

# Motivation

- GPUs are an interesting platform
  - Harness commodity hardware
  - Low cost
  - Easy to develop
  - Forwards compatible
- How competitive are GPUs?
  - Vs. other commodity hardware
  - Vs. FPGAs
  - Testcase: MWA Correlator
- Recall Nbody with GRAPE

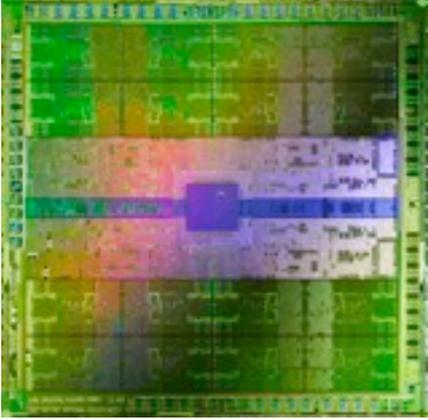
# Previous Work

- Harris *et al* (2008)
  - GPU for X-engine only - no consideration for integrated system
- Wayth *et al* (2009)
  - Consider GPU for both F-engine and X-engine
  - Developed for the 32 station MWA prototype
  - Not a scalable solution
- van Nieuwpoort *et al* (2009)
  - Compared X-engine performance across range of platforms
  - Claim: GPU implementations suffer from memory bottlenecks
  - Conclusion: BG/P and Cell optimal
- All of the above had low percentage of peak performance for GPUs



# MWA Correlator

- Array located in Western Australian Outback
- Total power budget ~ 50 kW
- Tasked with detecting EOR signal
- 512 stations x 2 polarizations = 1024 inputs
- total bandwidth = 31 MHz
- X-engine requires
  - $1.6 \times 10^{13}$  CMACs = 128 TFLOPS
- How many GPUs required?



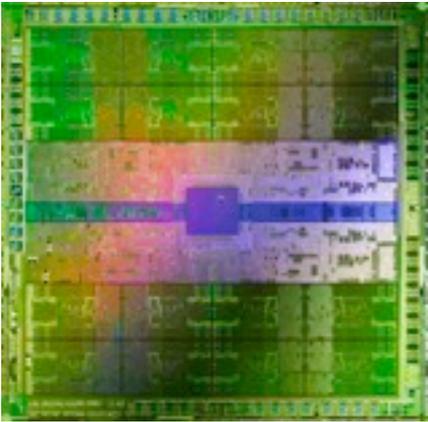
# Fermi Architecture

## GeForce GTX 480

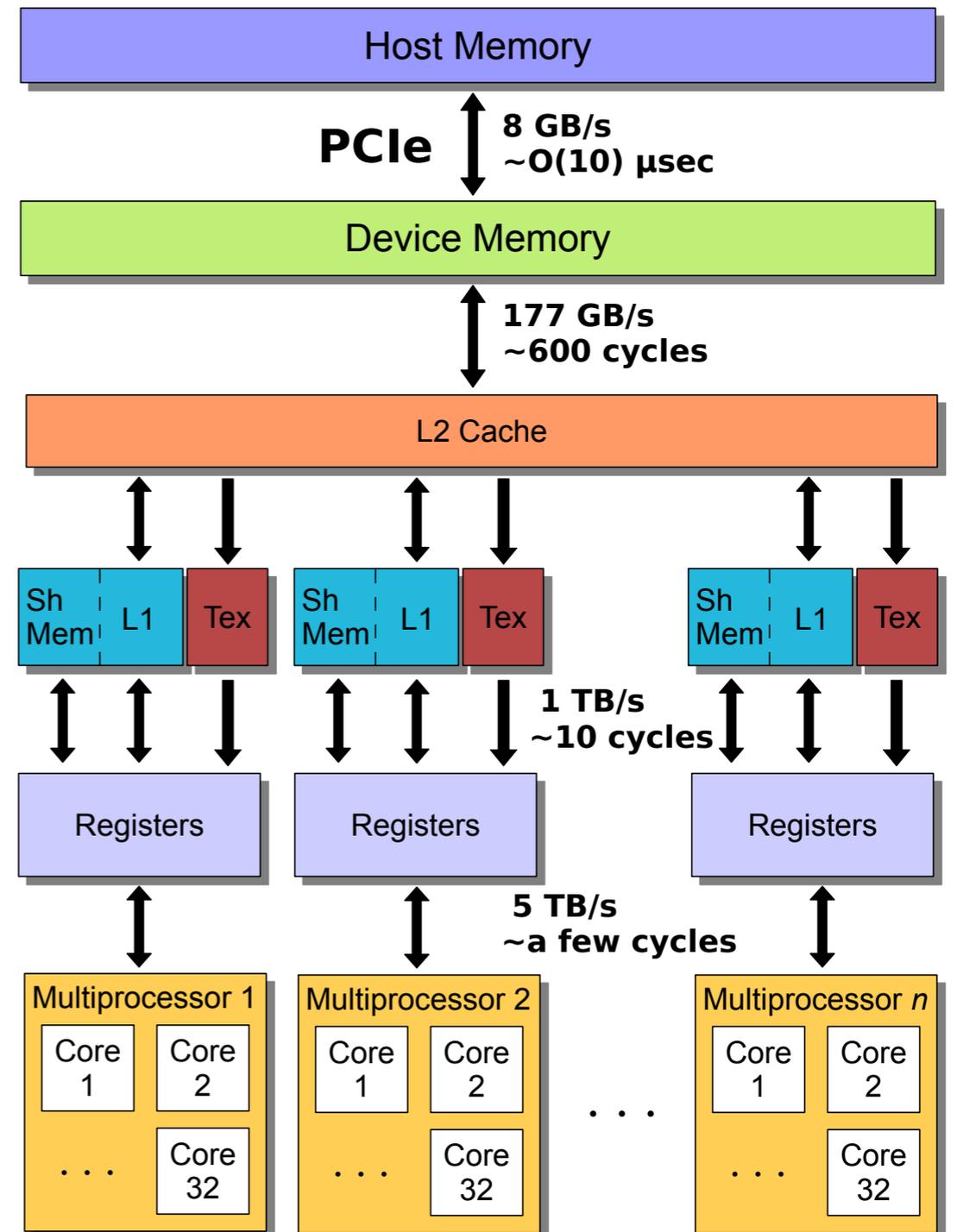
- 480 processing cores
- 1.345 TFLOPS SP
- 177 GB/s memory bandwidth
- Power consumption 250 Watts
- Easy to program
- \$300 (as of yesterday)



# Fermi Architecture



- High memory bandwidth
- BUT high flop:byte ratio 7.6:1
- Need high arithmetic intensity to keep GPU busy
- Use memory hierarchy
  - Registers
  - Shared memory
  - Device memory
  - PCIe bus

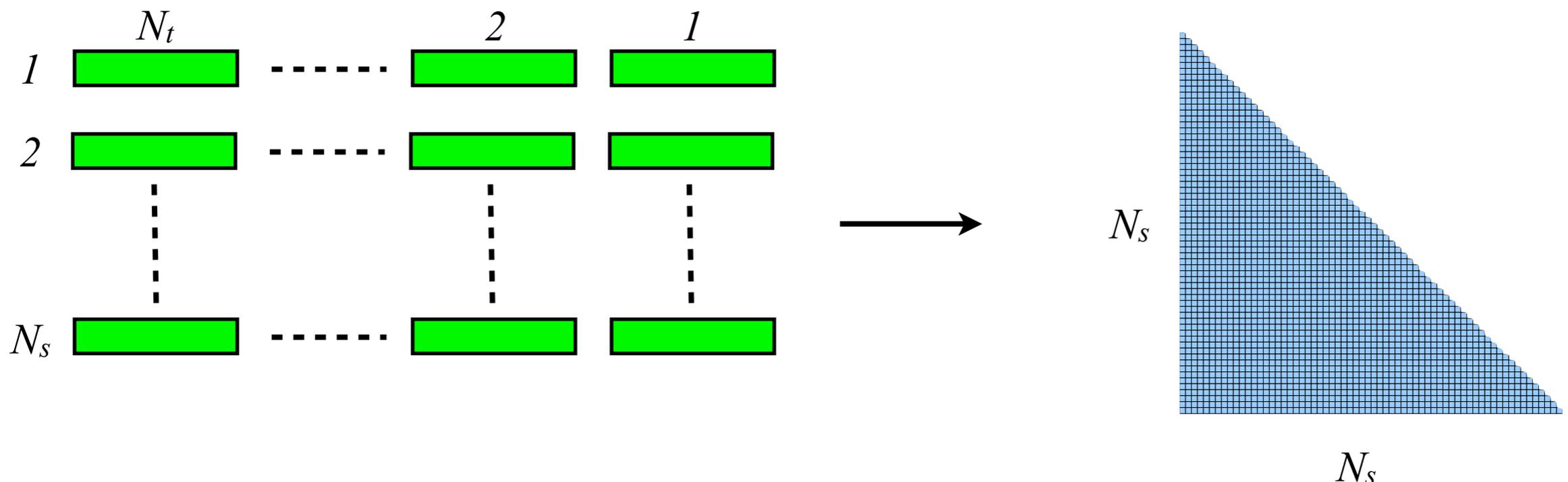


# The X-engine

- Mathematically, just the sum of a series of vector outer products

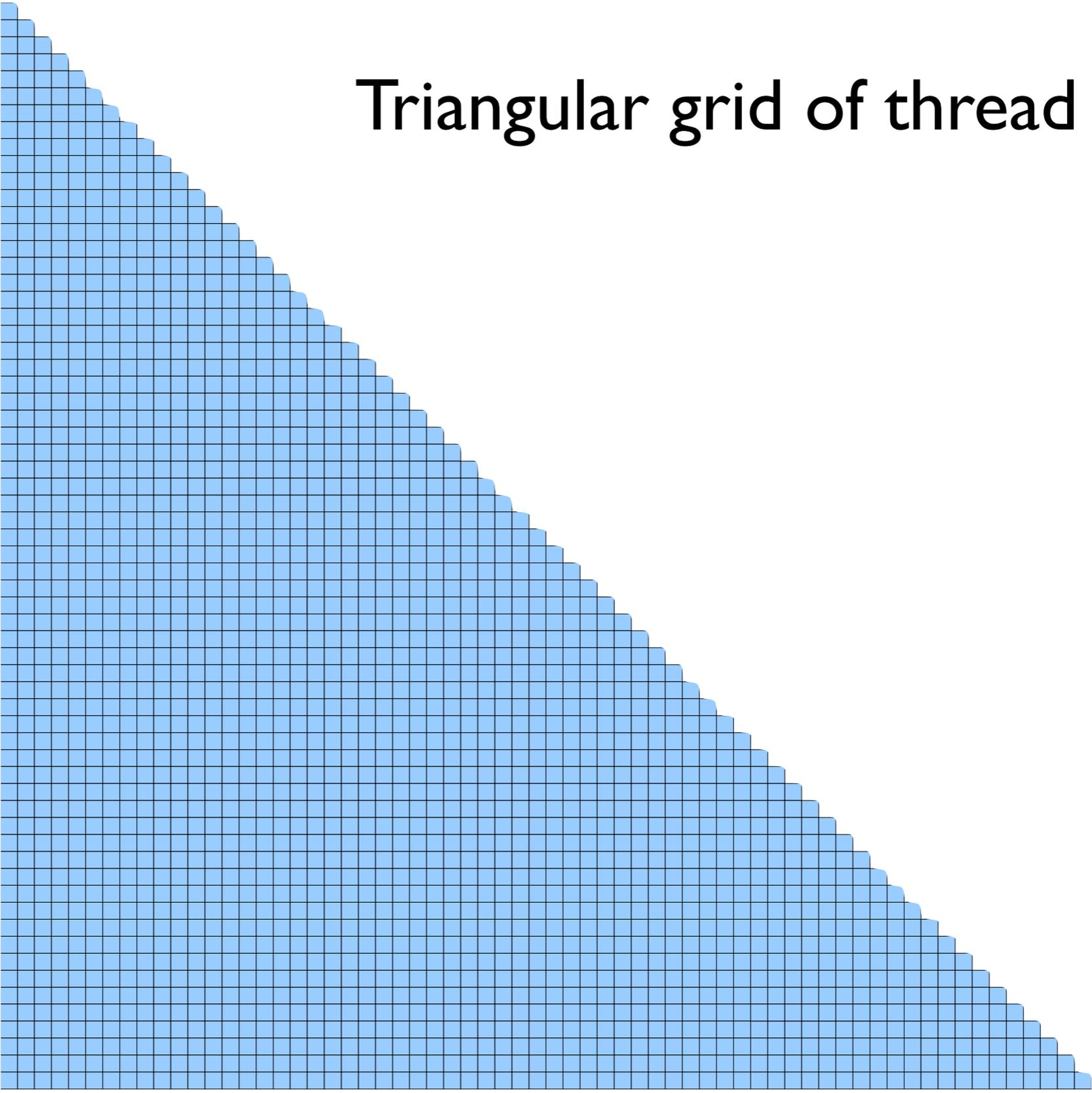
$$S_{ij}(\nu) = \sum_{t=1}^{N_t} X_i(\nu) X_j(\nu)^\dagger \quad \text{FLOPS} = \frac{1}{2} 8B(2N_s)(2N_s + 1)$$

- $N_t=B/T$  is the integration length, e.g.  $\sim 100000$
- Matrix is Hermitian - just calculate lower triangular elements



# Triangular grid of thread blocks

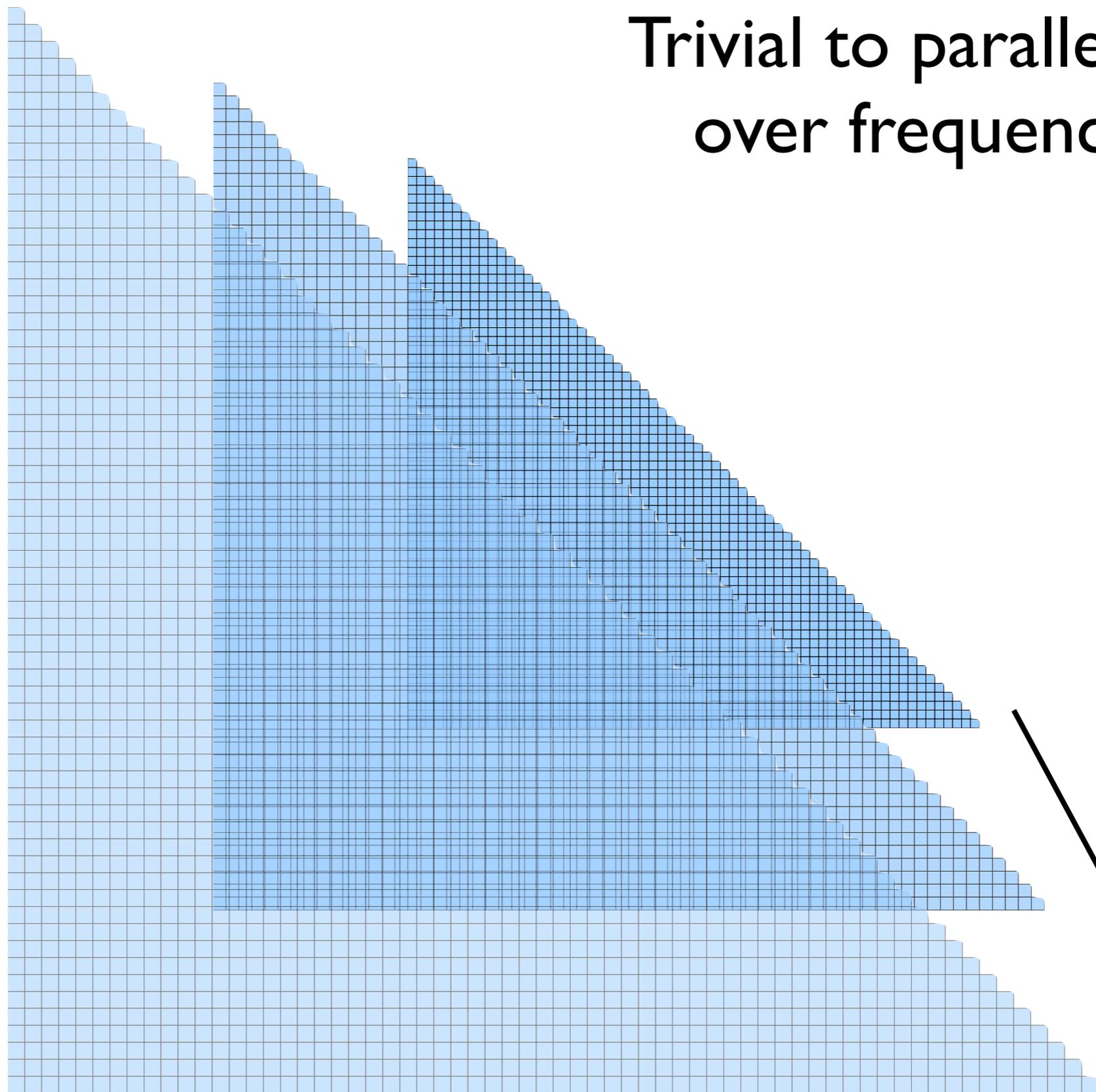
Station



Station

Trivial to parallelize  
over frequency

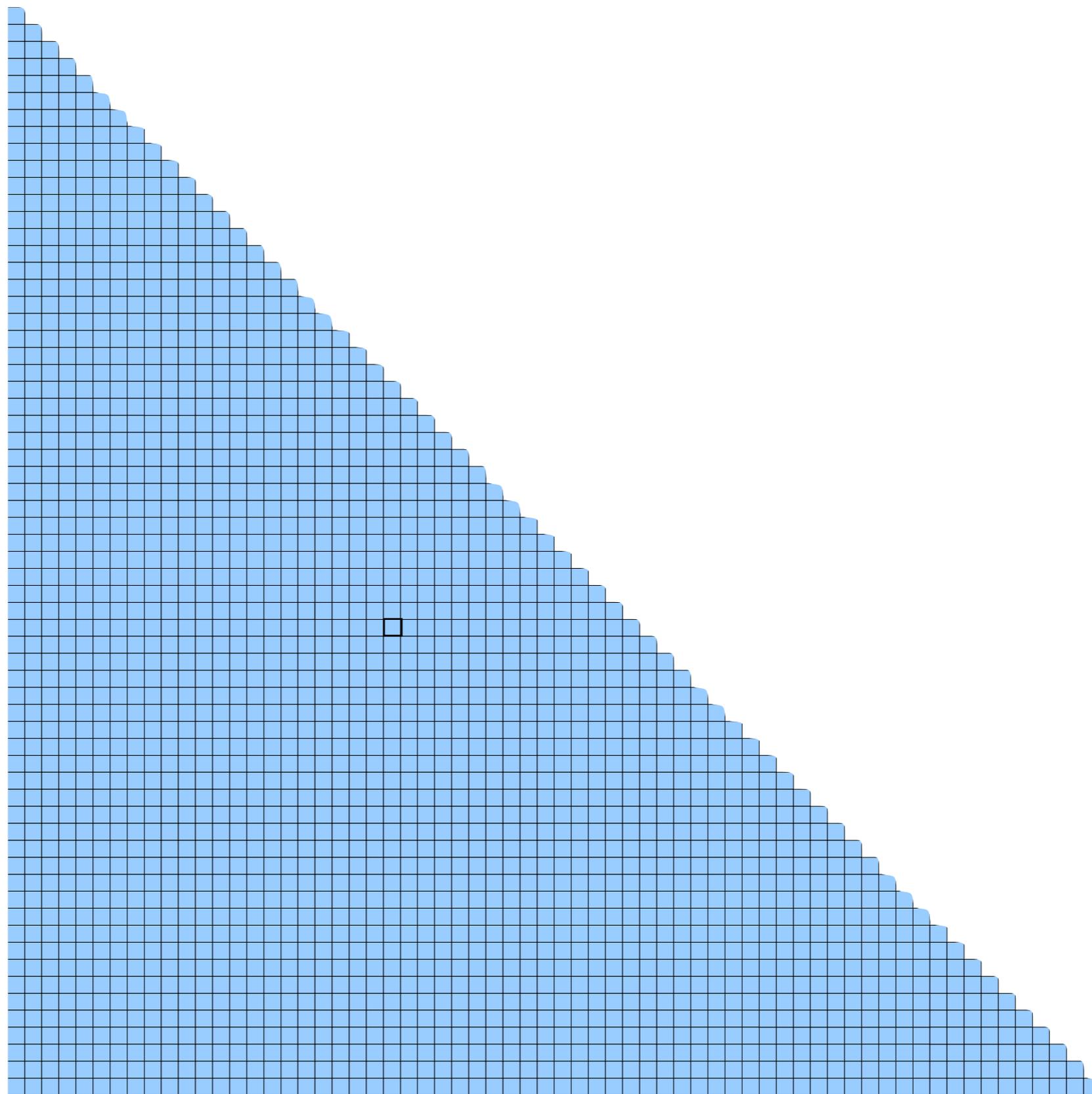
Station



$N_c$

Station

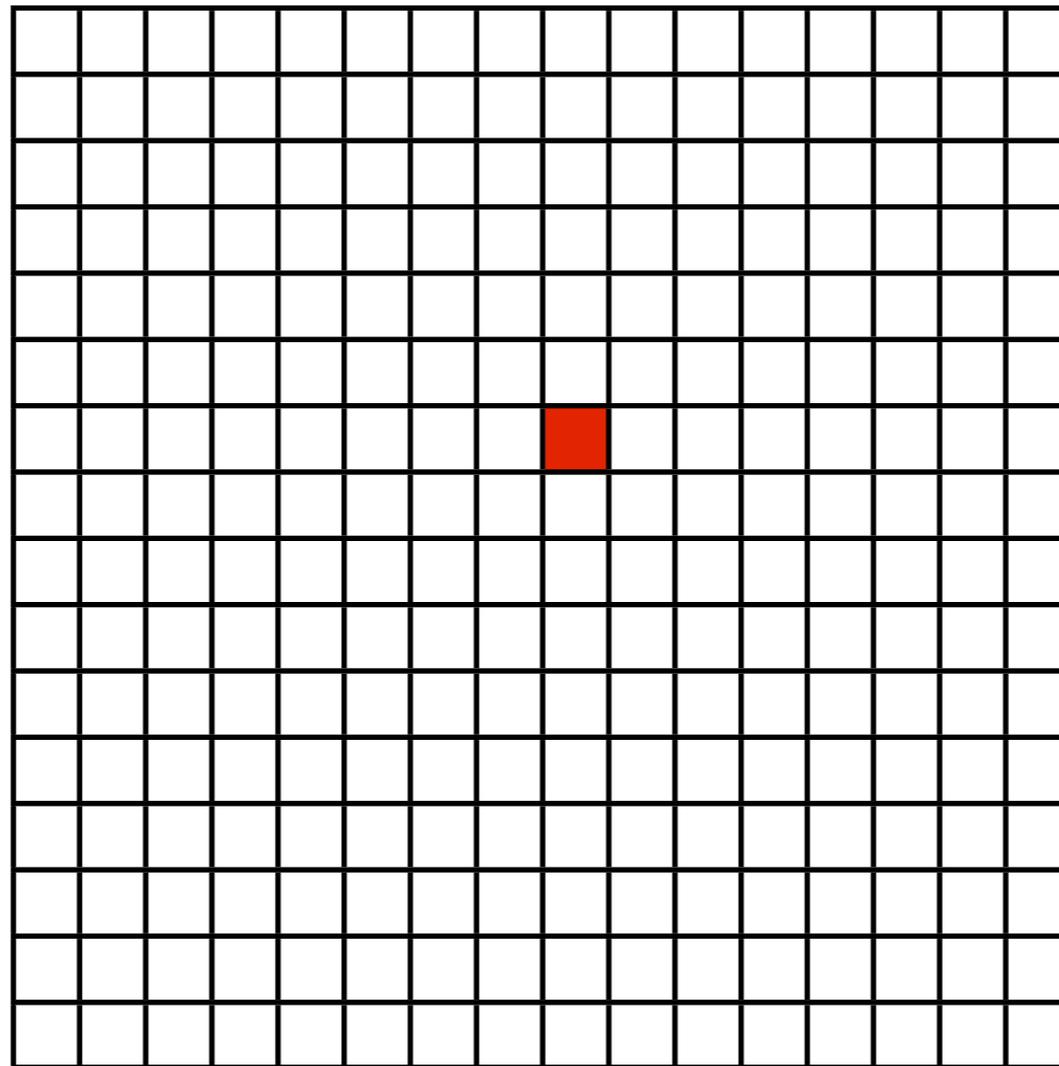
Station



Station

16x16 threads = 16x16 stations

Matrix elements stored in registers



## Registers

Each thread  
computes  
a 1x1 tile

flop/byte  
Algorithm: I  
Hardware: 7.6

## Cache

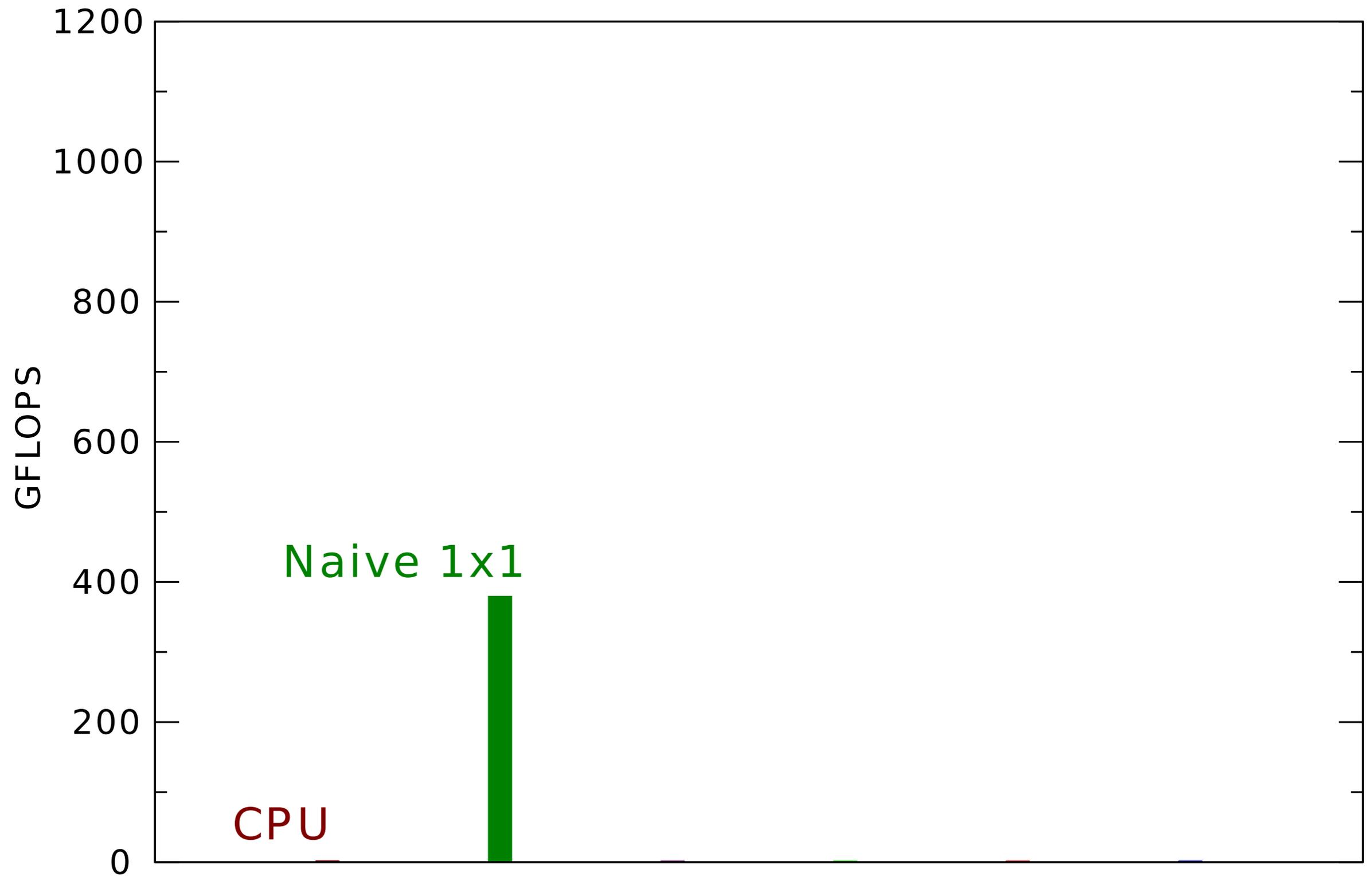
Data shared  
between threads  
using L1 cache

## Output

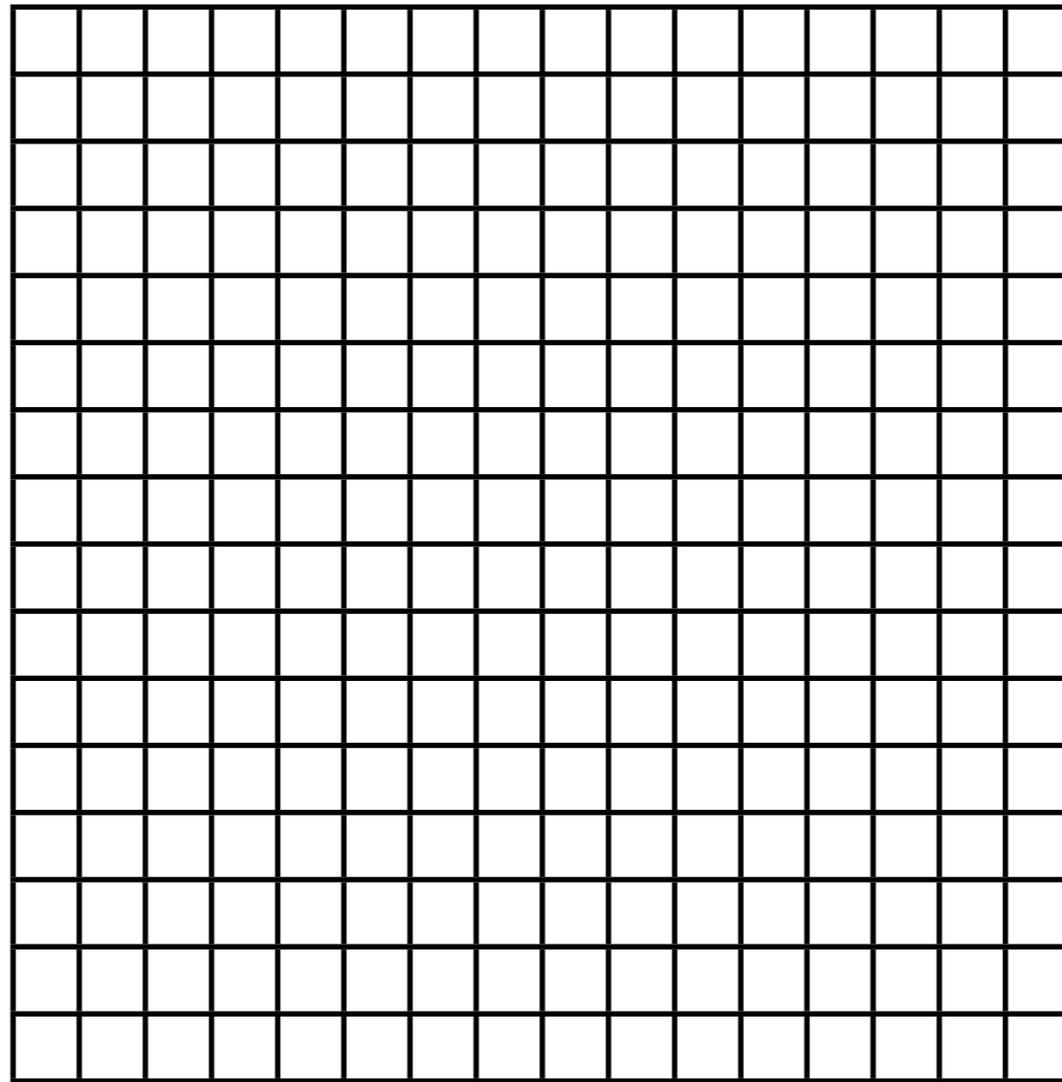
Each thread outputs its  
1x1 tile to memory

flop/byte  
Algorithm: I/N<sub>t</sub>  
Hardware: 7.6

# Performance for $N_s=512$ , $N_c=12$ , $N_t=1000$



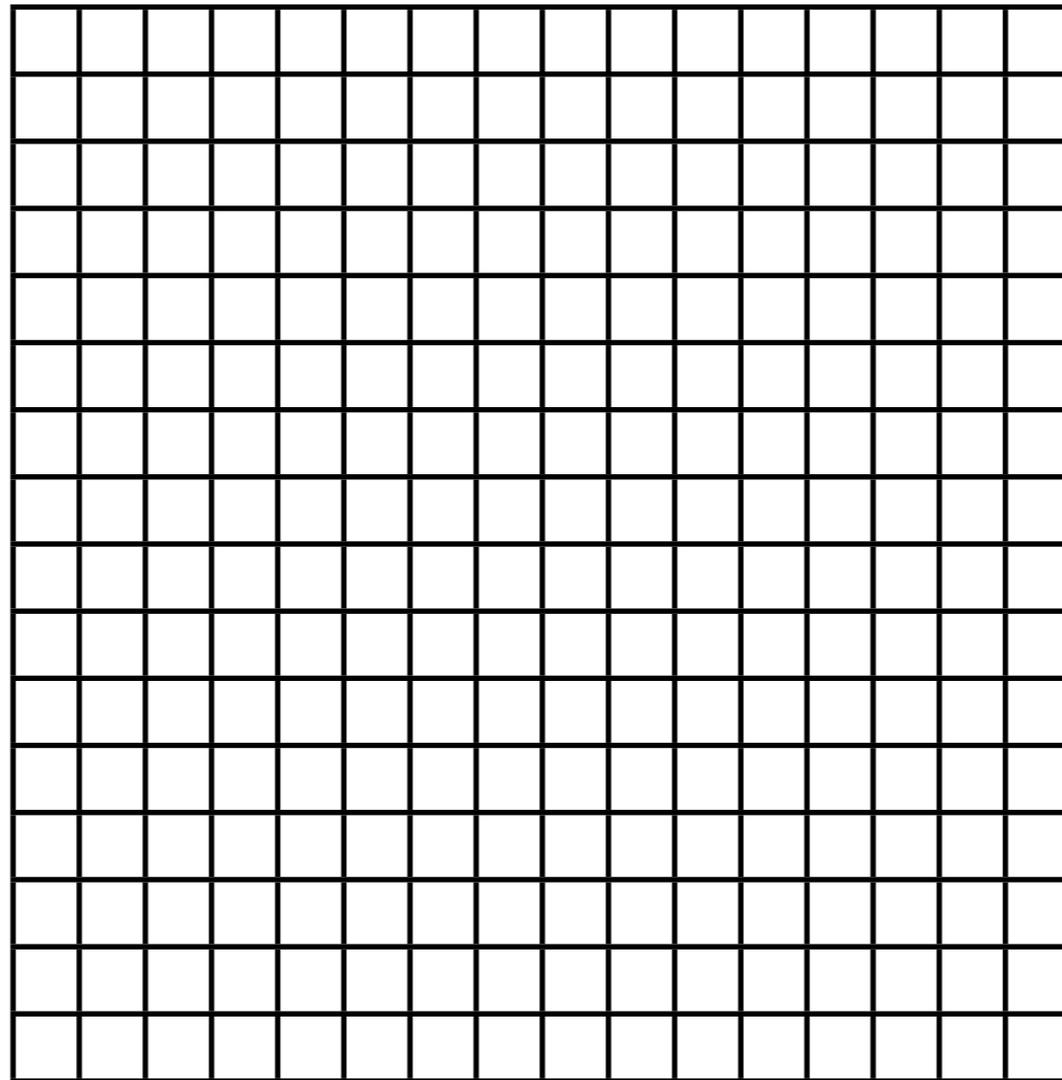
# Matrix elements stored in registers



shared memory



# Matrix elements stored in registers

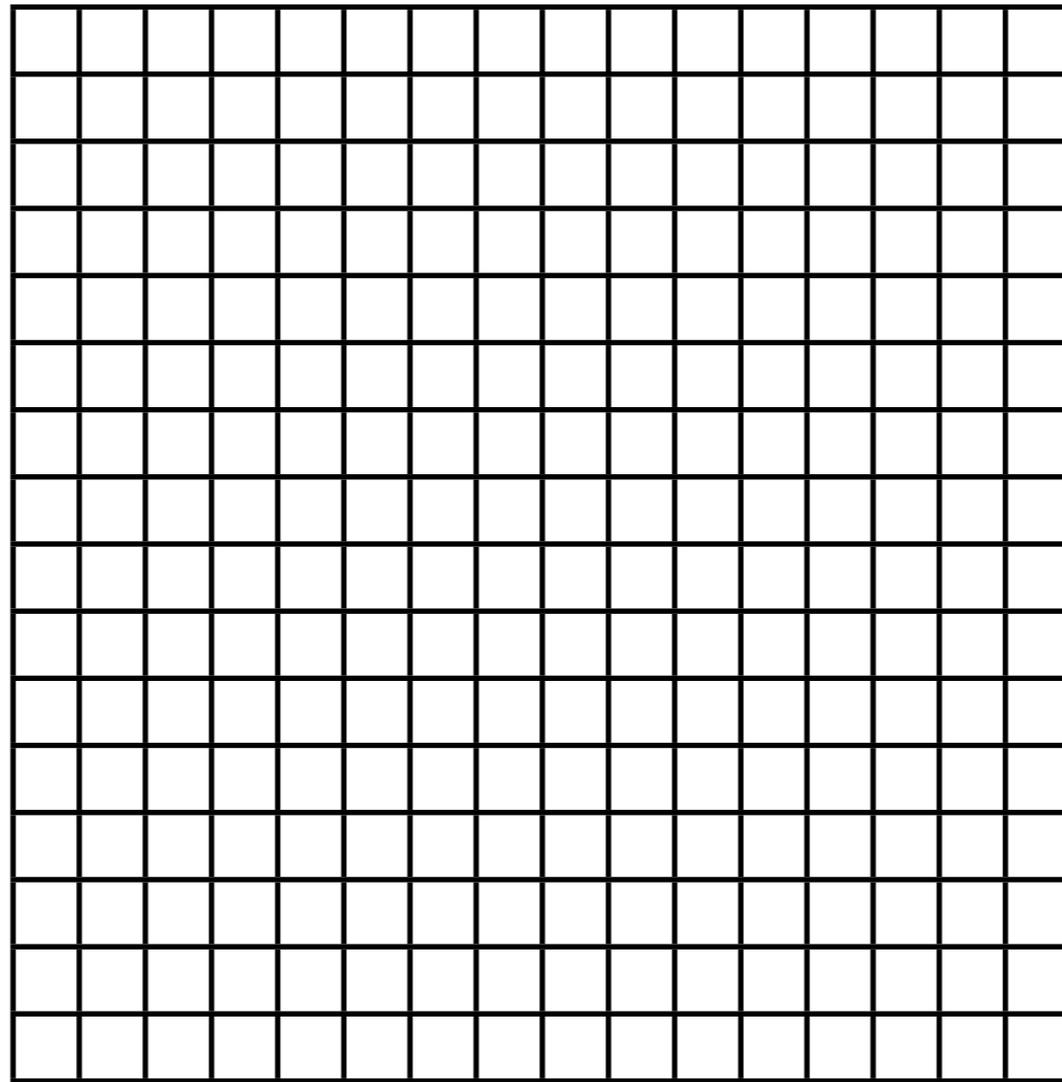


shared memory



1st warp reads row

# Matrix elements stored in registers

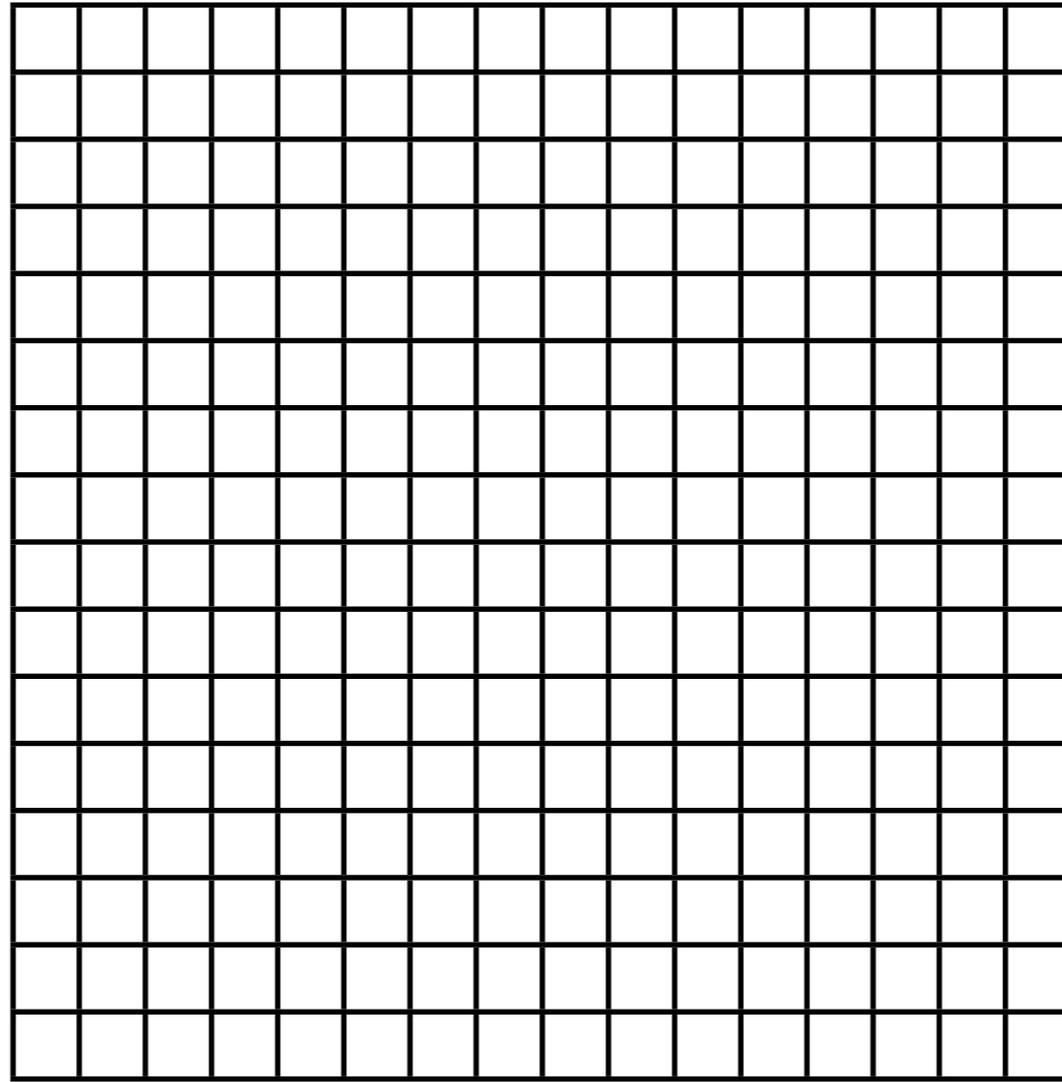


shared memory



# Matrix elements stored in registers

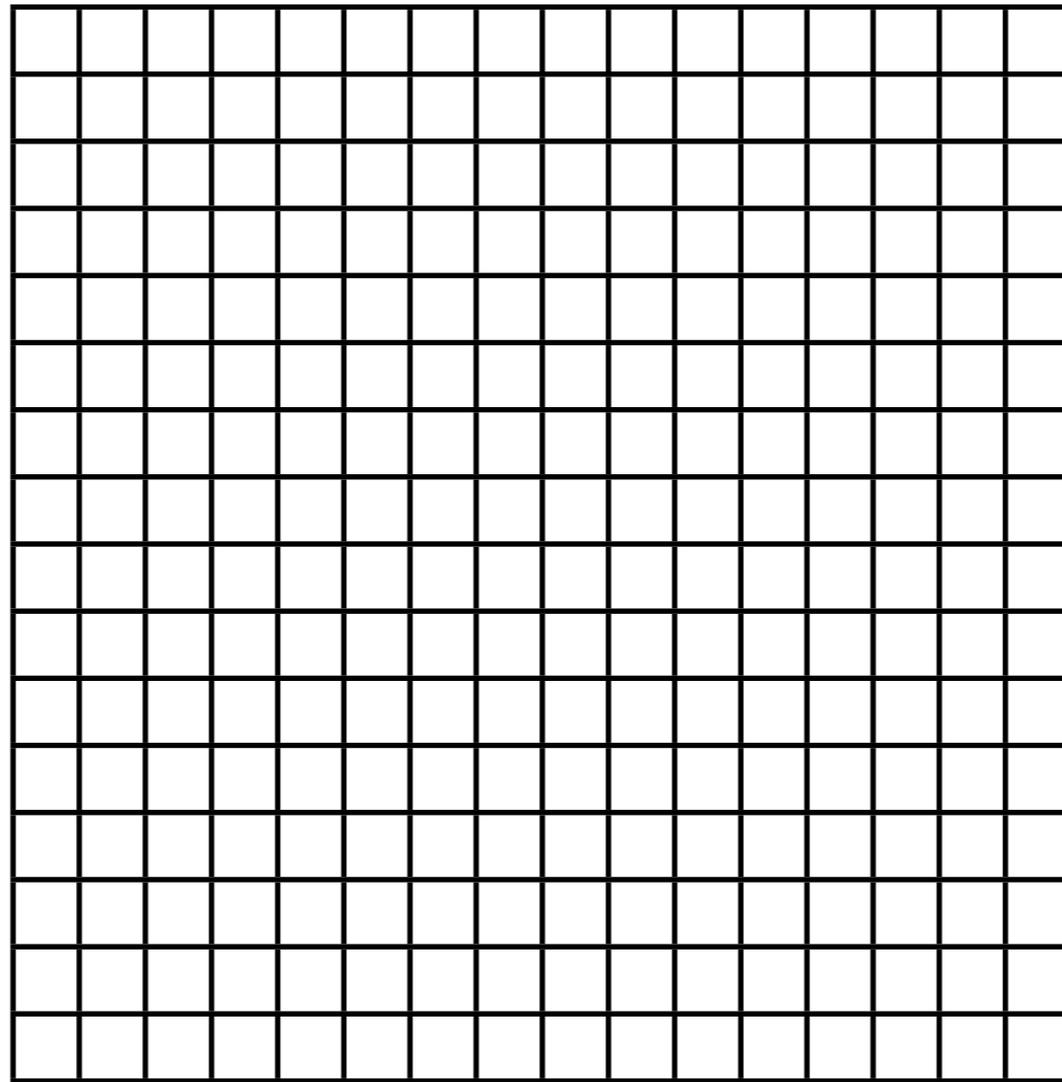
2nd warp reads column



shared memory

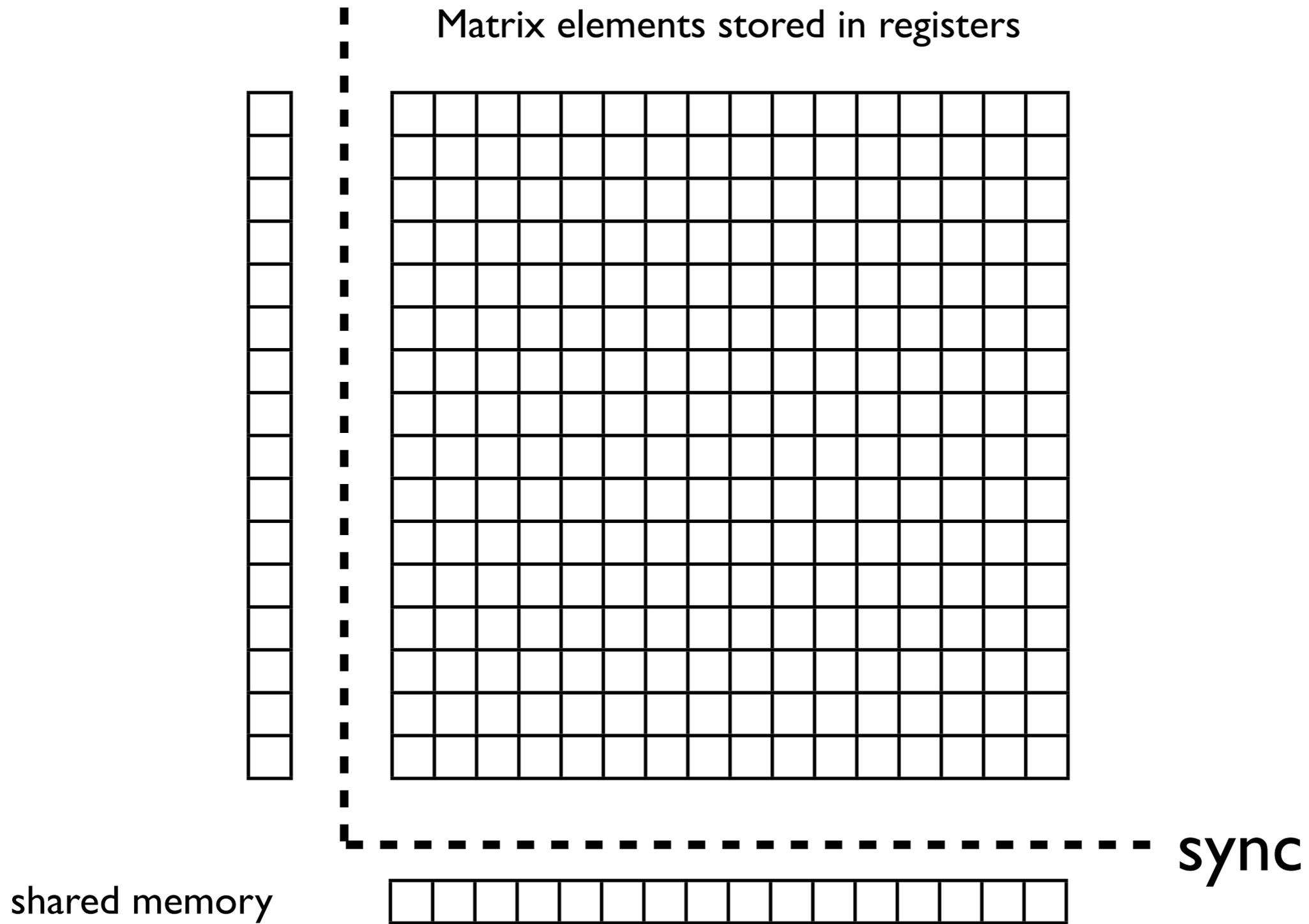


# Matrix elements stored in registers

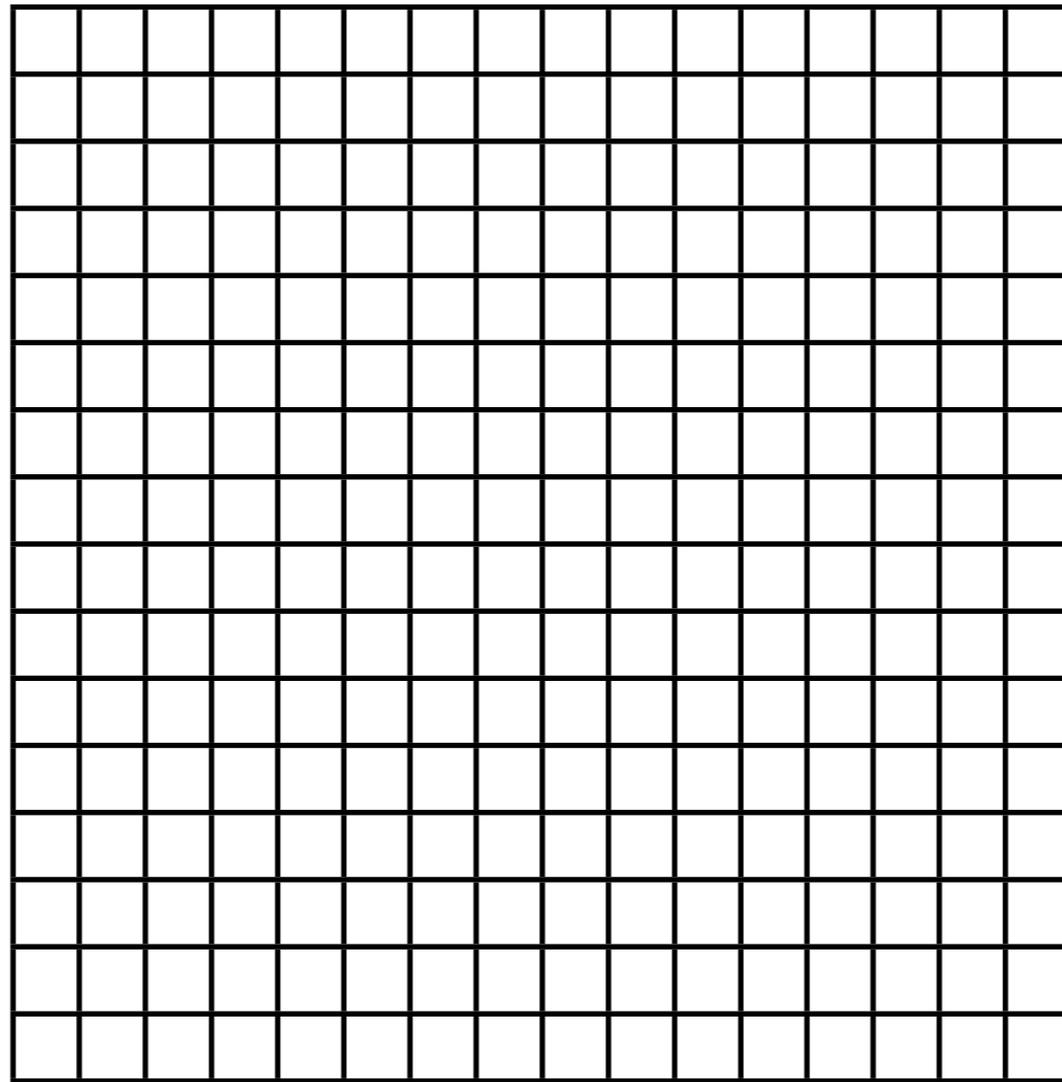


shared memory





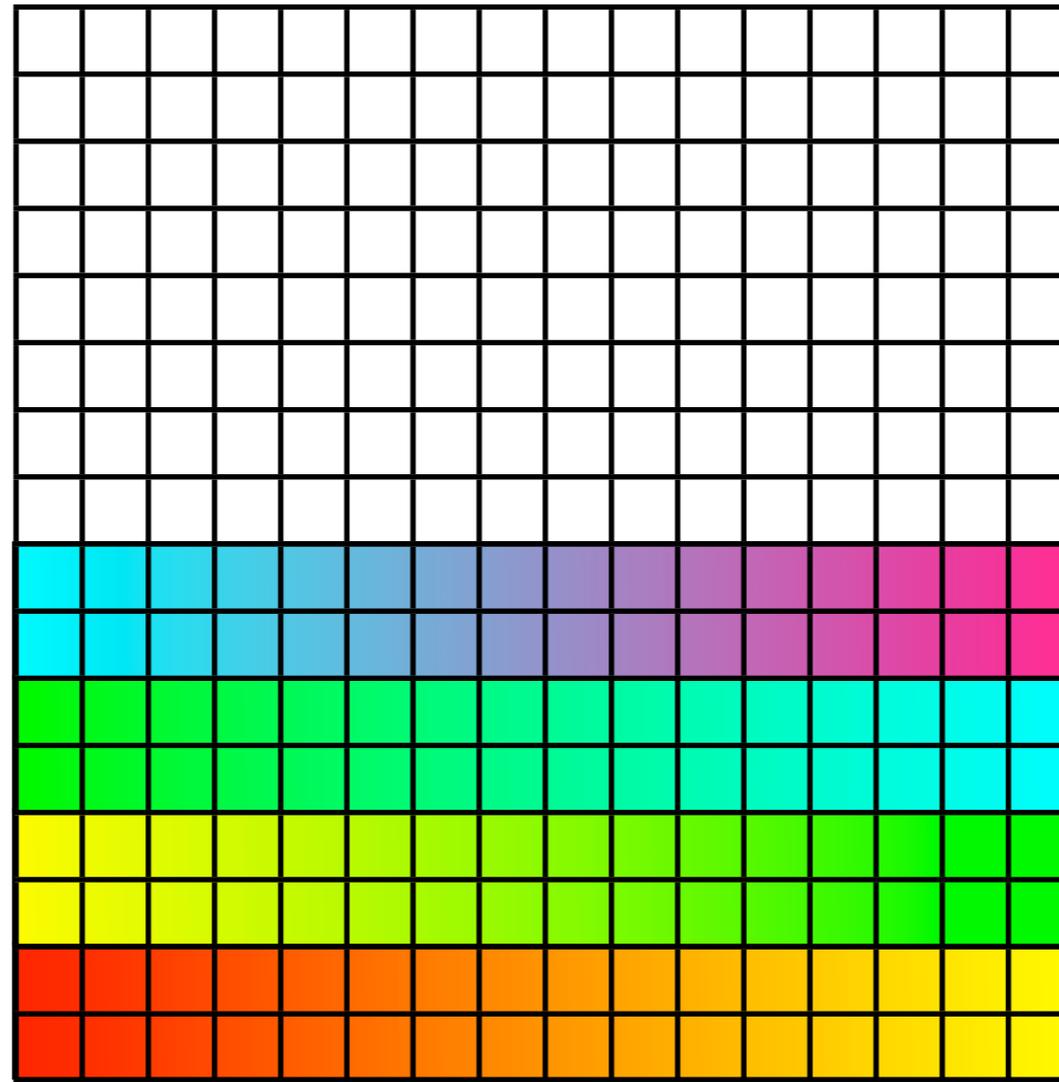
# Matrix elements stored in registers



shared memory



# Matrix elements stored in registers

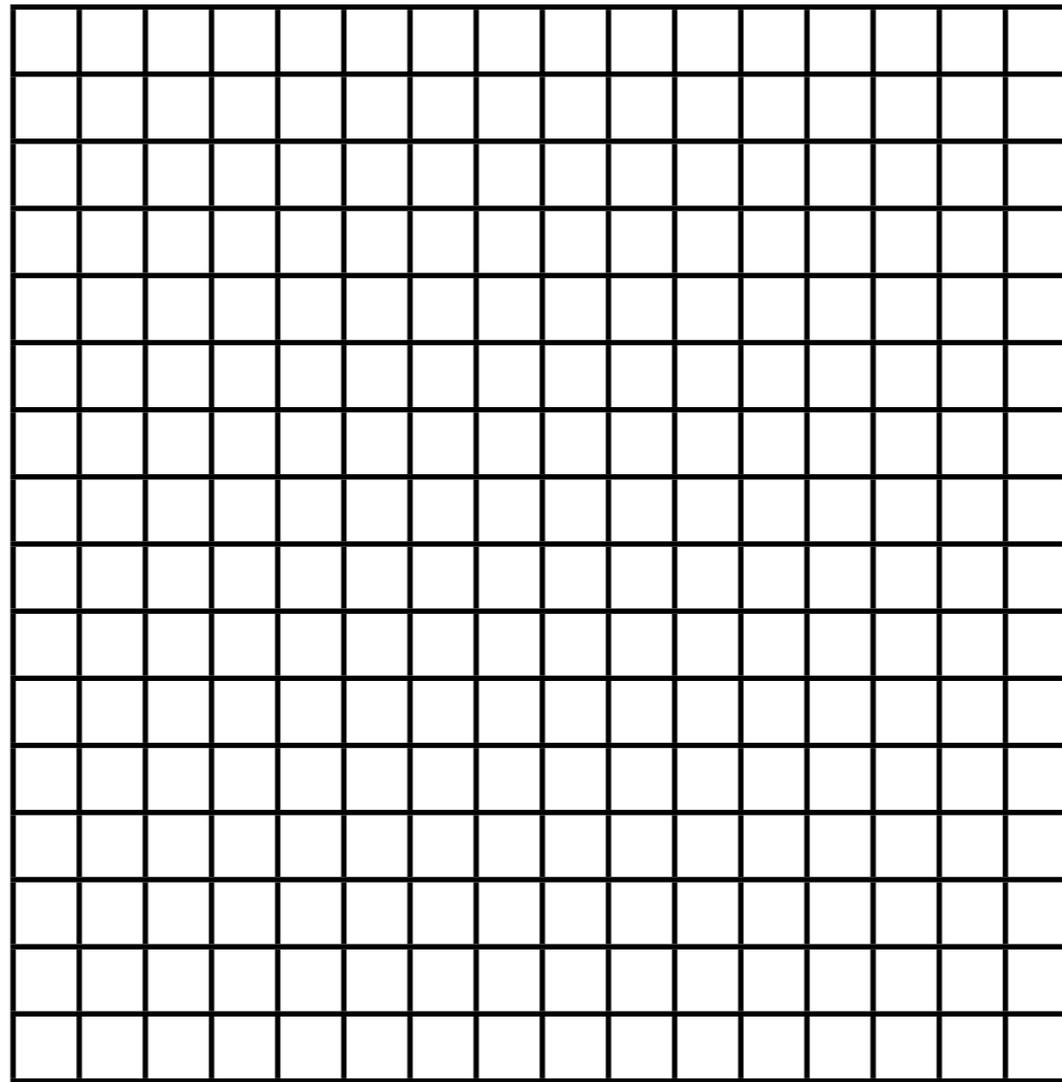


1st warp compute

shared memory



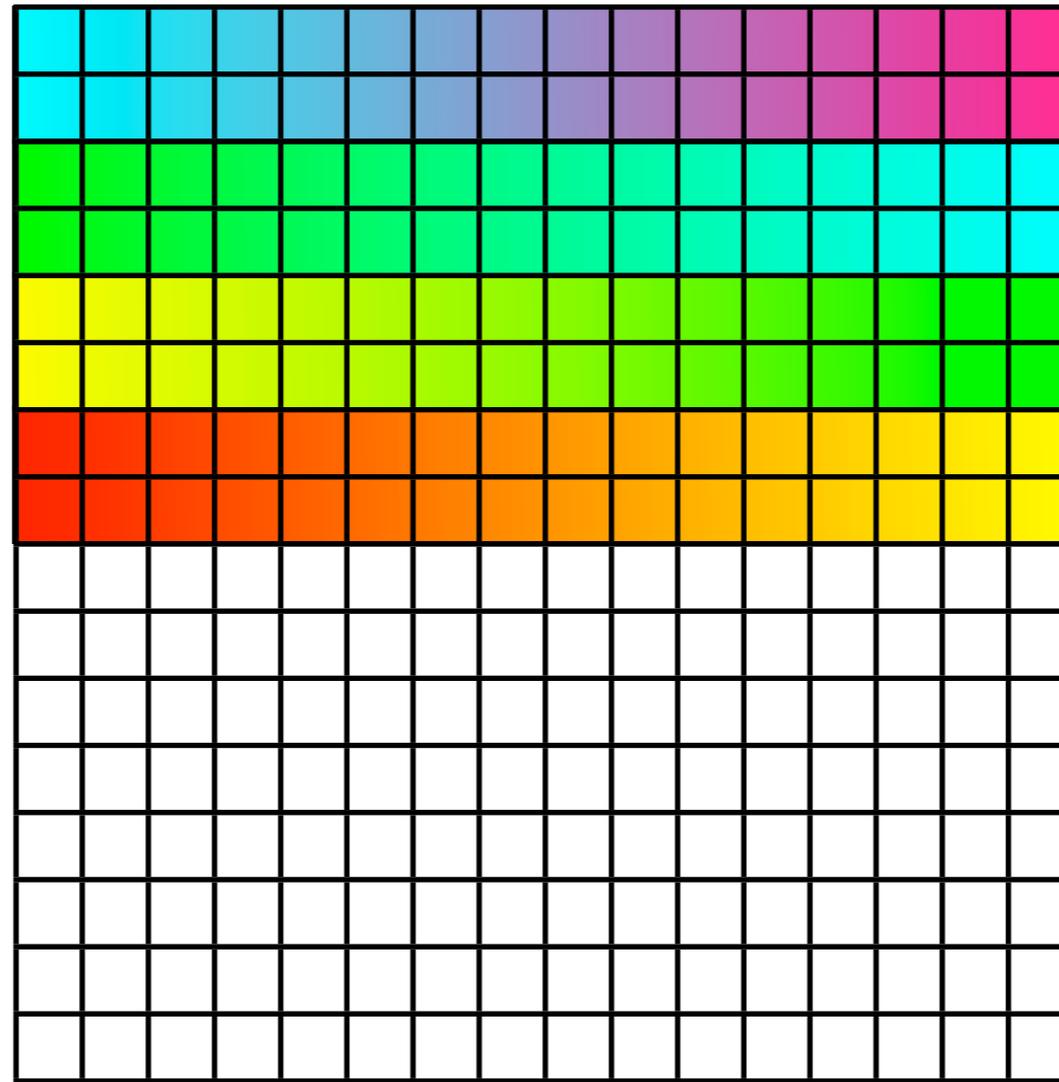
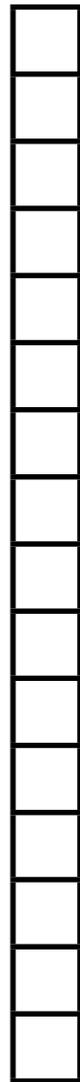
# Matrix elements stored in registers



shared memory



# Matrix elements stored in registers

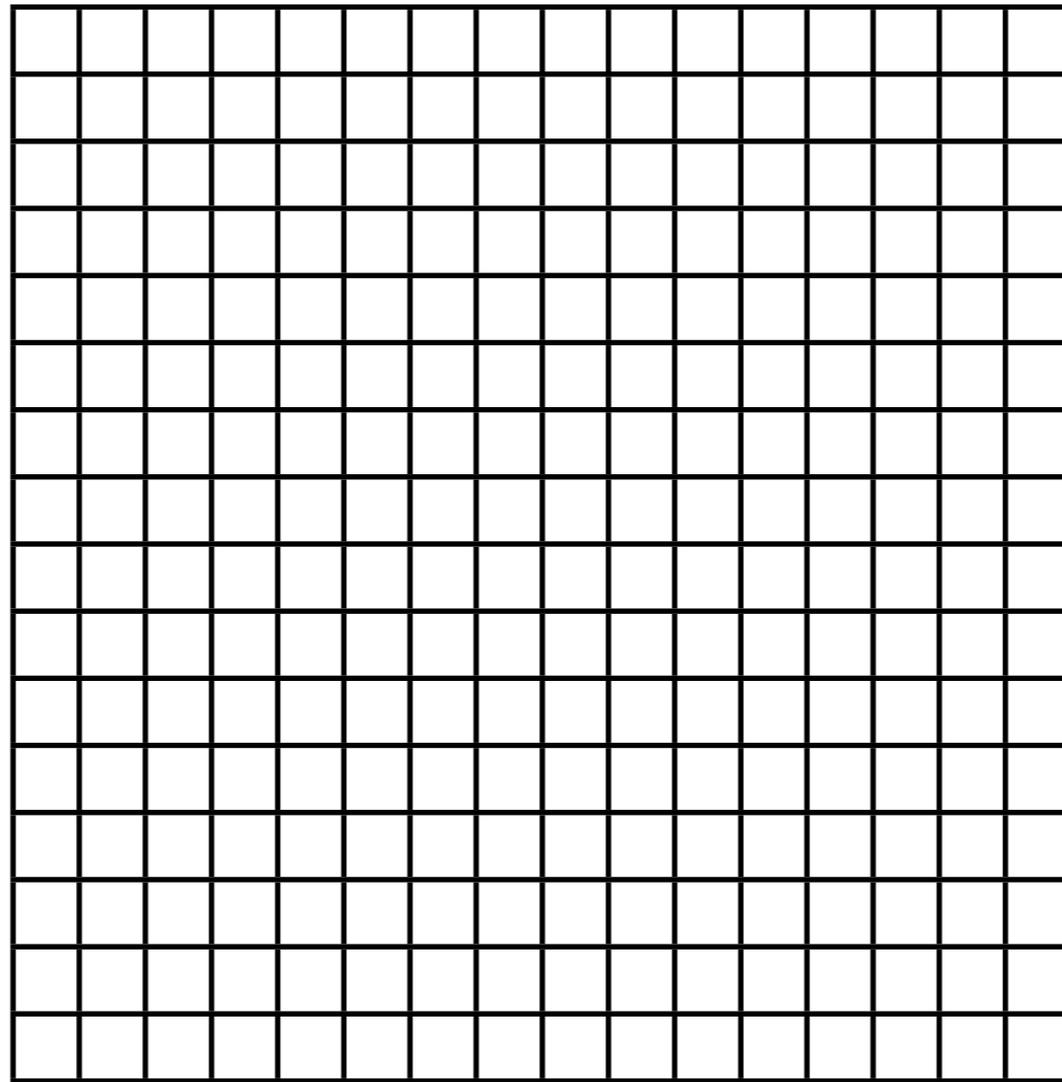


2nd warp compute

shared memory



# Matrix elements stored in registers



shared memory



## Registers

Each thread  
computes  
a 1x1 tile

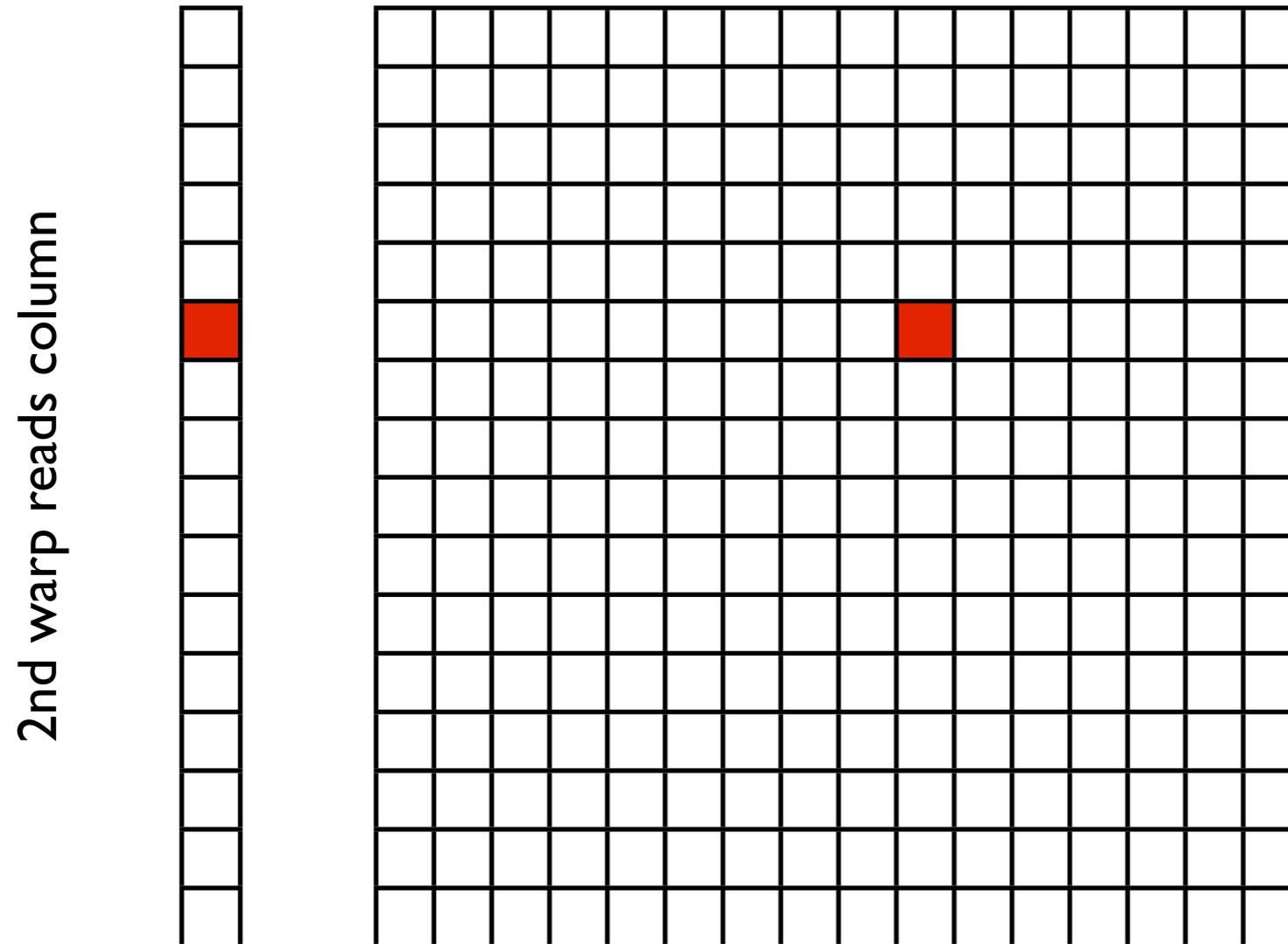
flop/byte  
Algorithm: 1  
Hardware: 1.5

## Shared Memory

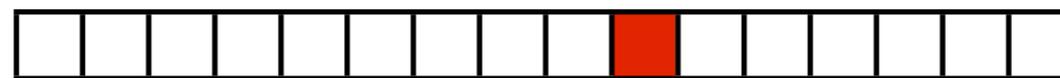
Each thread  
block  
loads a 16x16 tile

flop/byte  
Algorithm: 16  
Hardware: 7.6

Matrix elements stored in registers

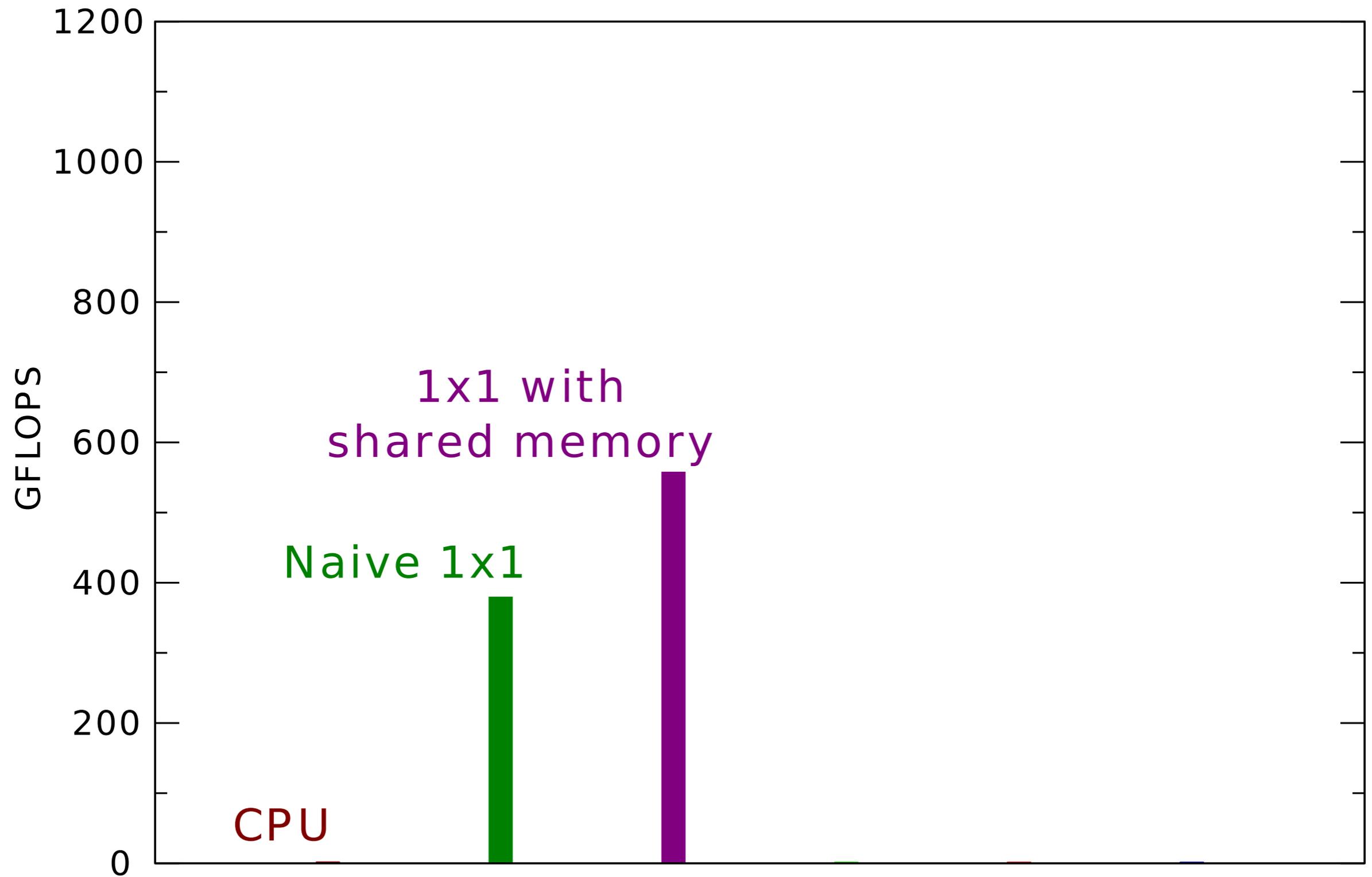


shared memory



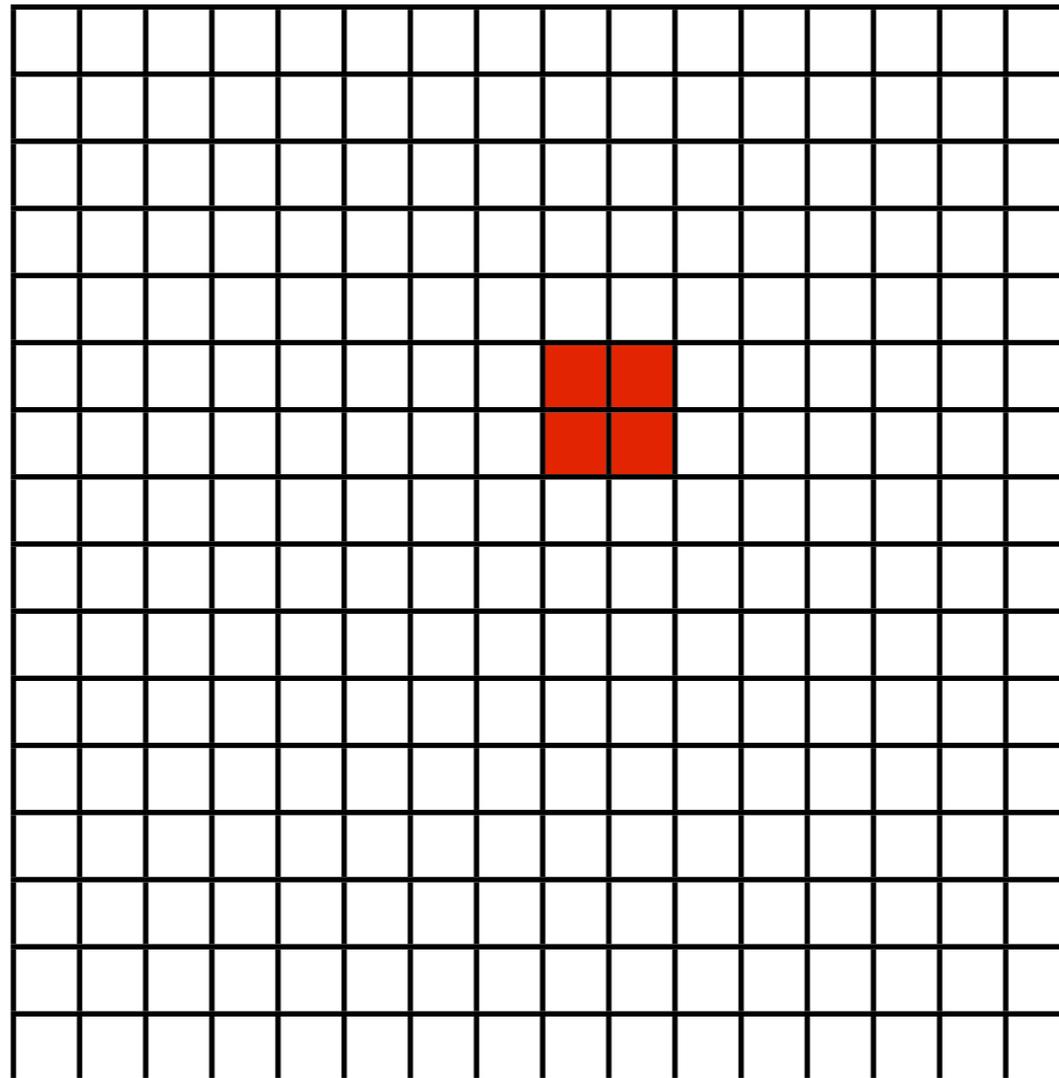
1st warp reads row

# Performance for $N_s=512$ , $N_c=12$ , $N_t=1000$



8x8 threads = 16x16 stations

Matrix elements stored in registers



## Registers

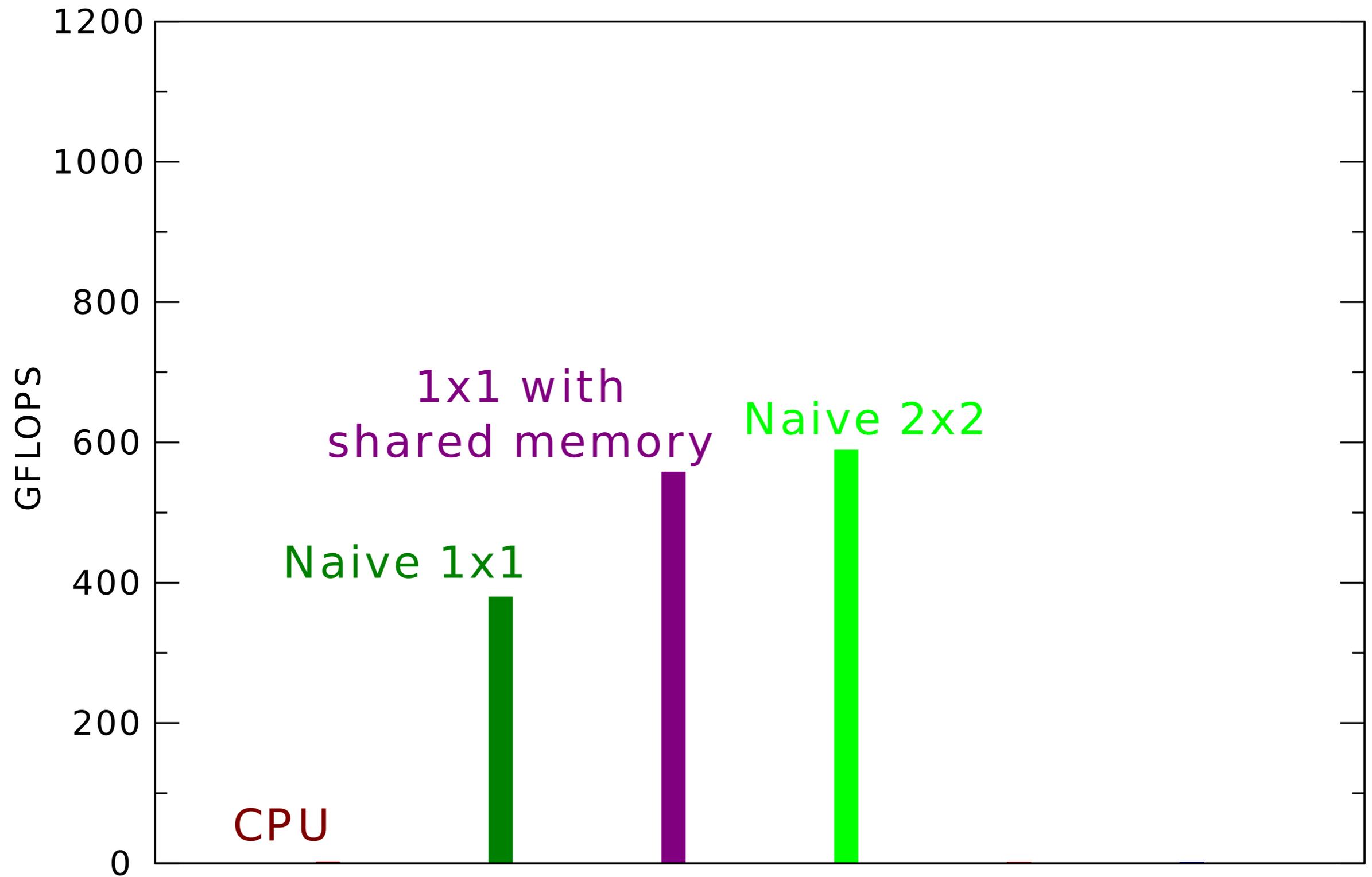
Each thread  
computes  
a 2x2 tile

flop/byte  
Algorithm: 2  
Hardware: 7.6

## Cache

Data shared  
between threads  
using L1 cache

# Performance for $N_s=512$ , $N_c=12$ , $N_t=1000$



## Registers

Each thread  
computes  
a  $1 \times 1$  tile

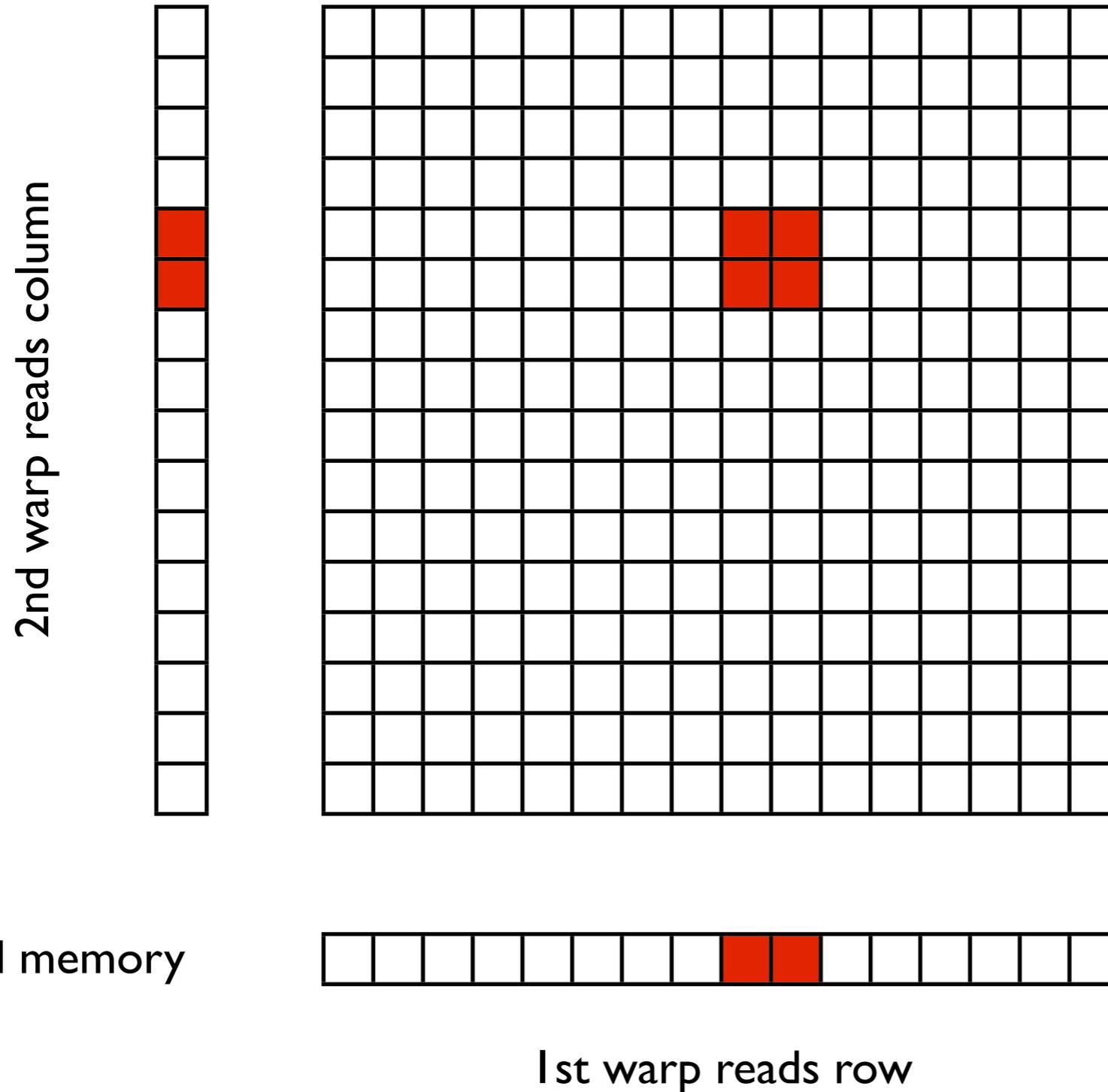
flop/byte  
Algorithm: 2  
Hardware: 1.5

## Shared Memory

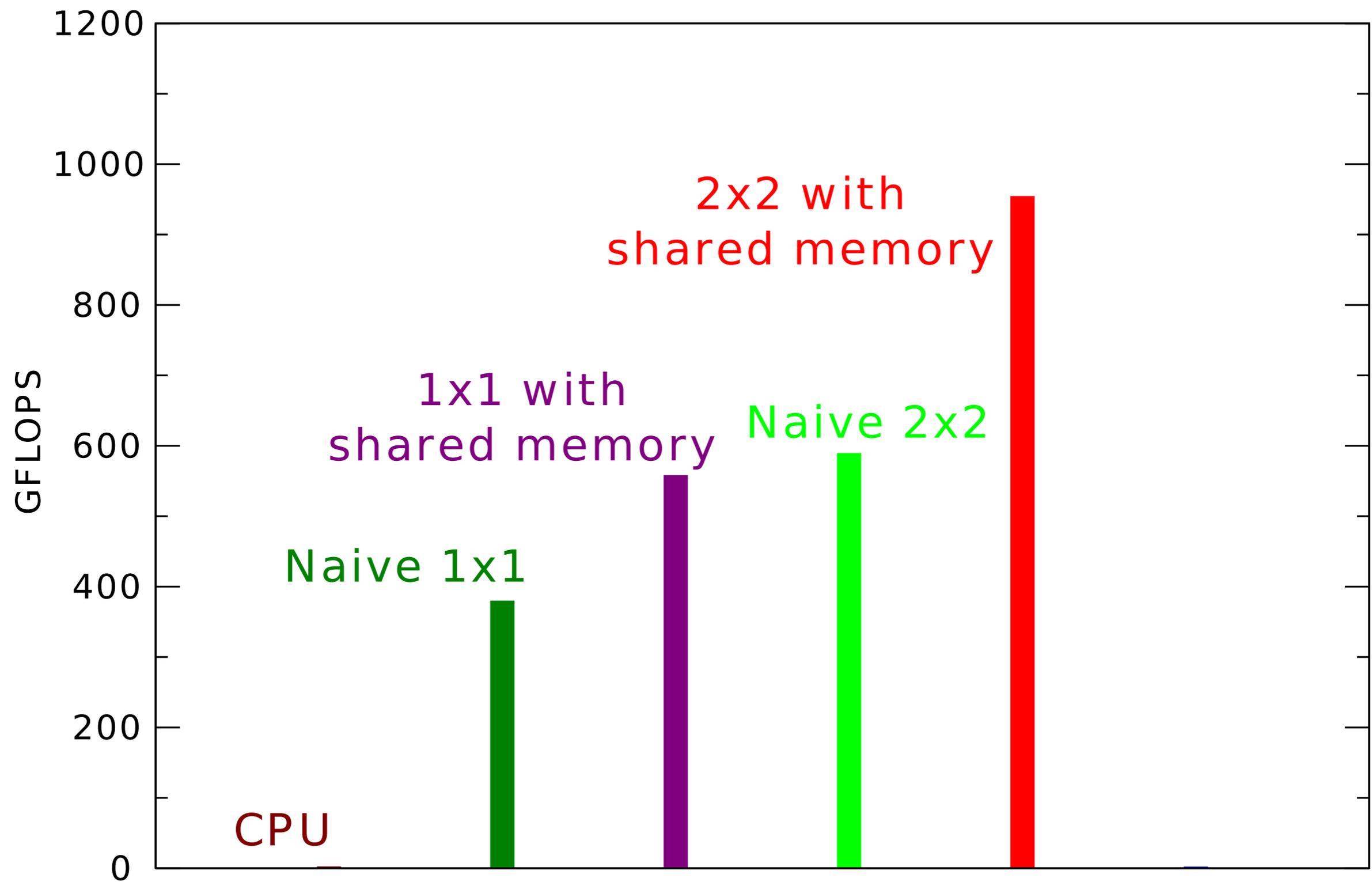
Each thread  
block  
loads a  $16 \times 16$  tile

flop/byte  
Algorithm: 16  
Hardware: 7.6

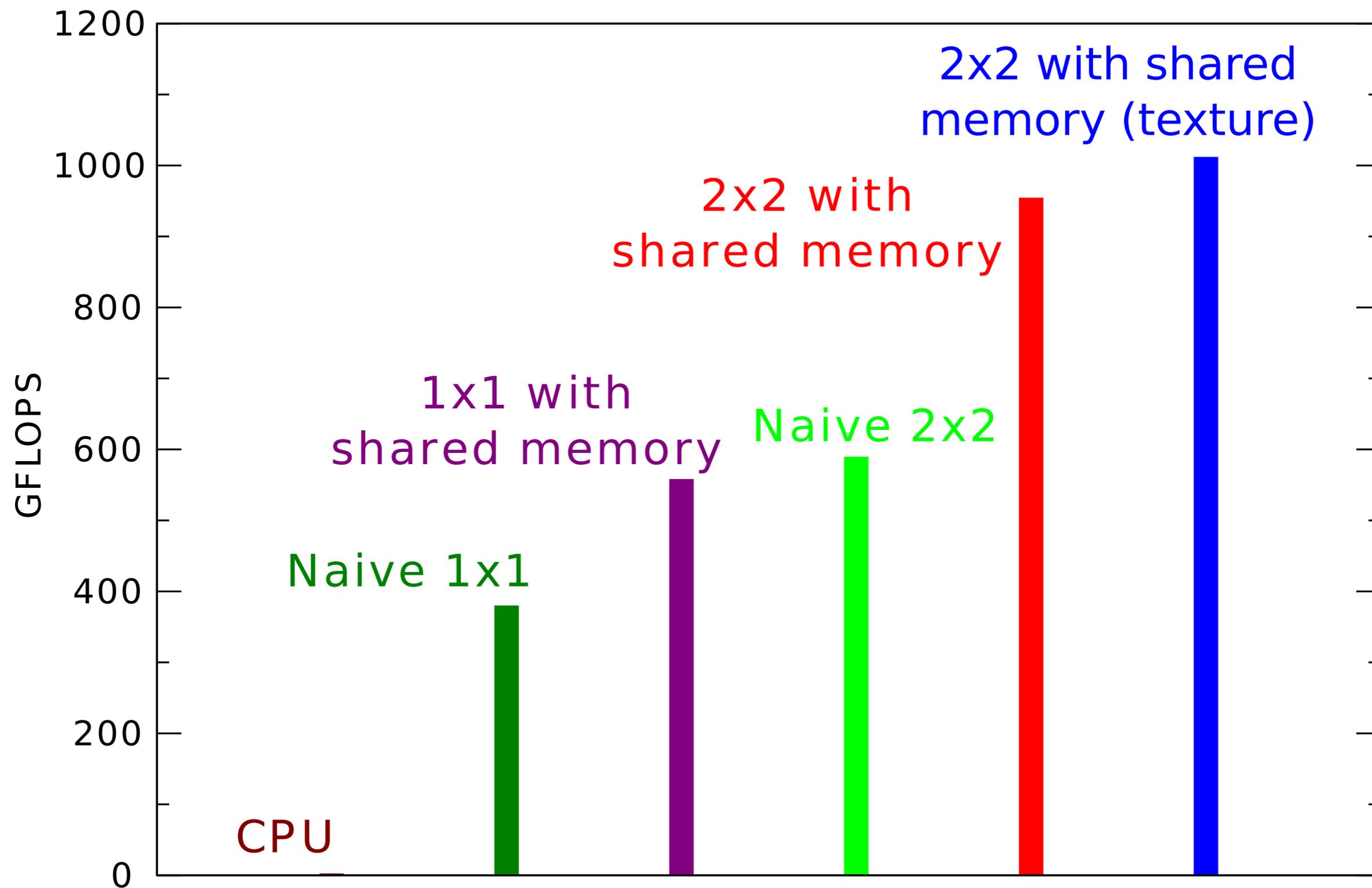
Matrix elements stored in registers



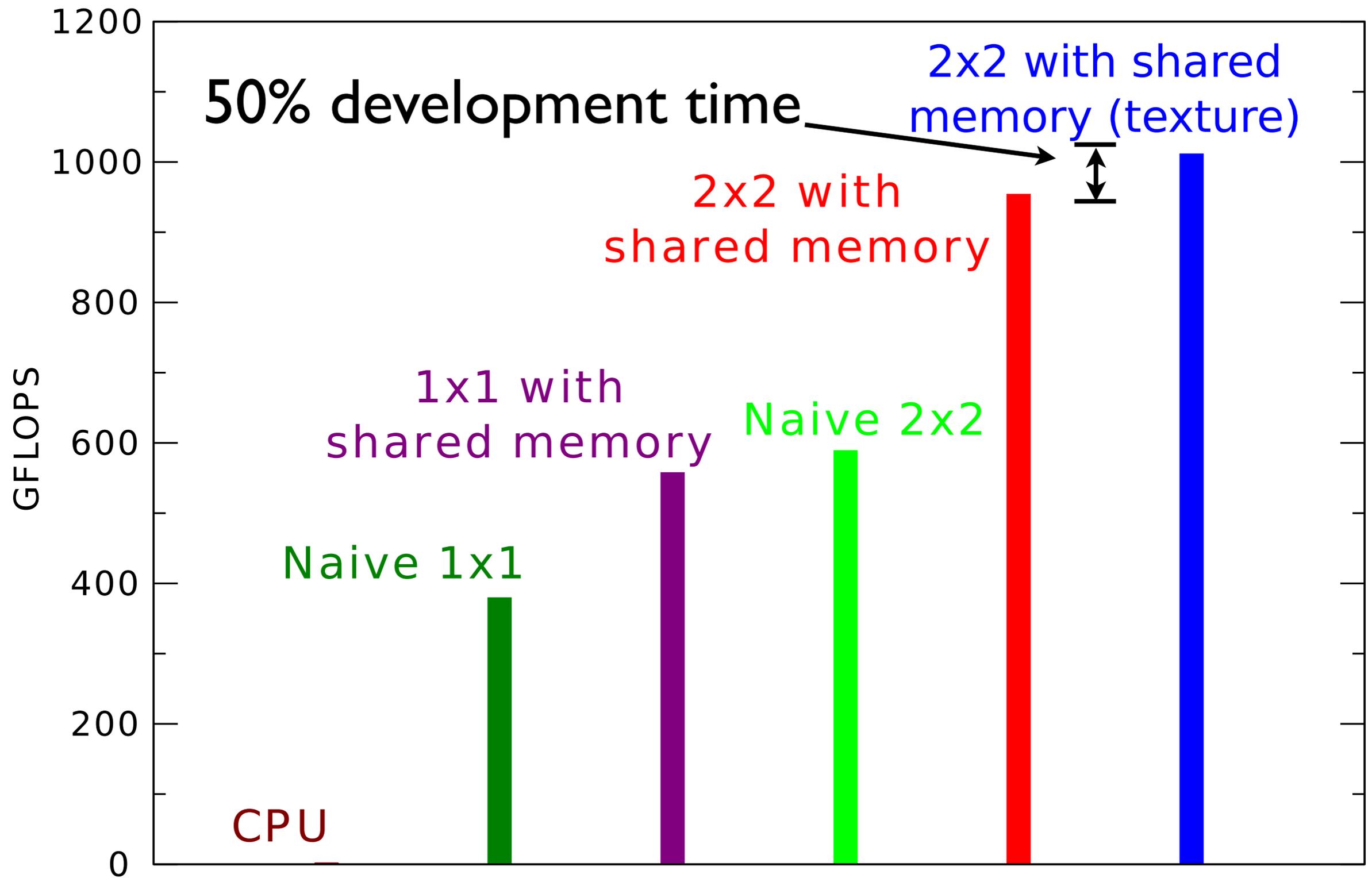
# Performance for $N_s=512$ , $N_c=12$ , $N_t=1000$

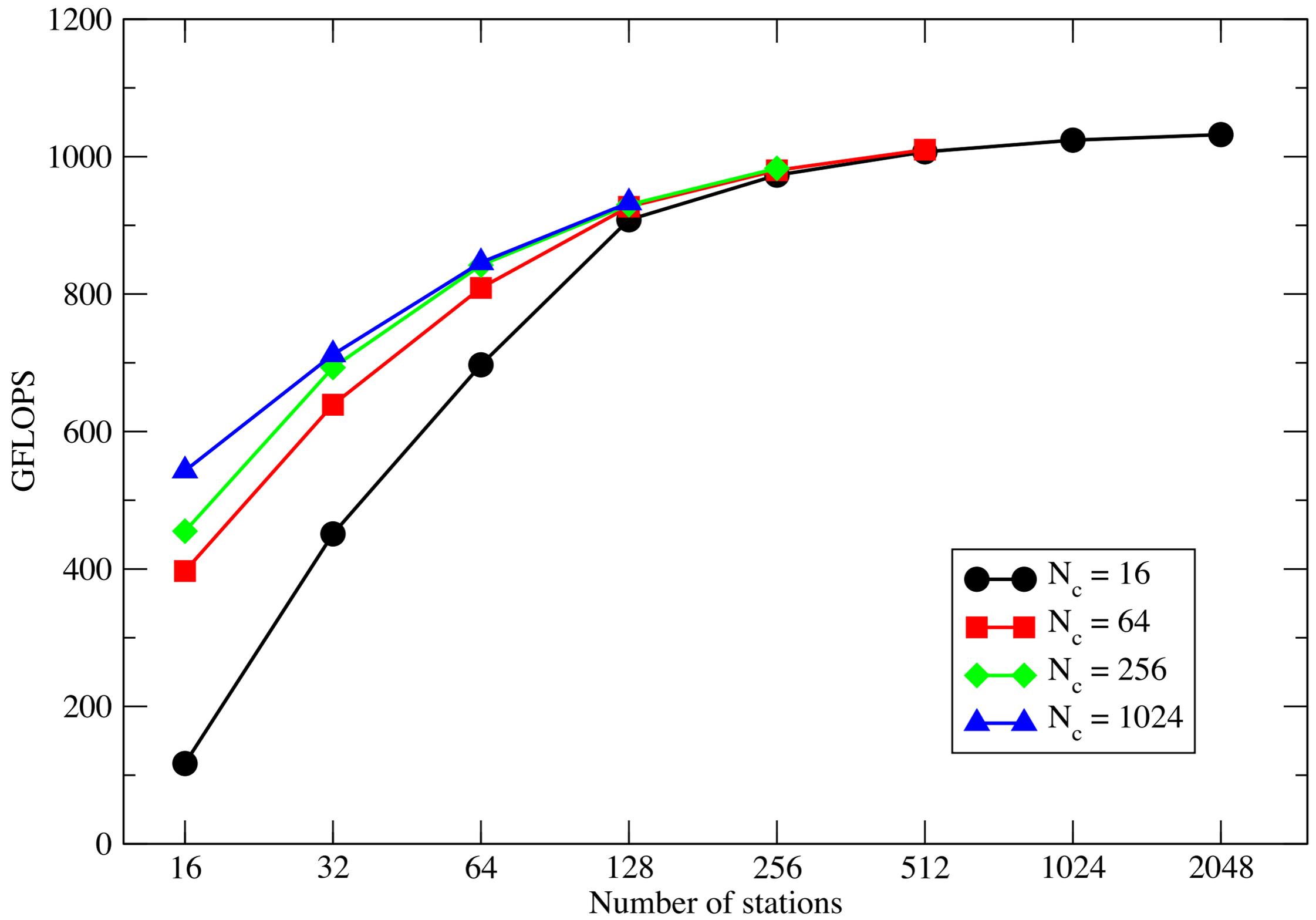


# Performance for $N_s=512$ , $N_c=12$ , $N_t=1000$

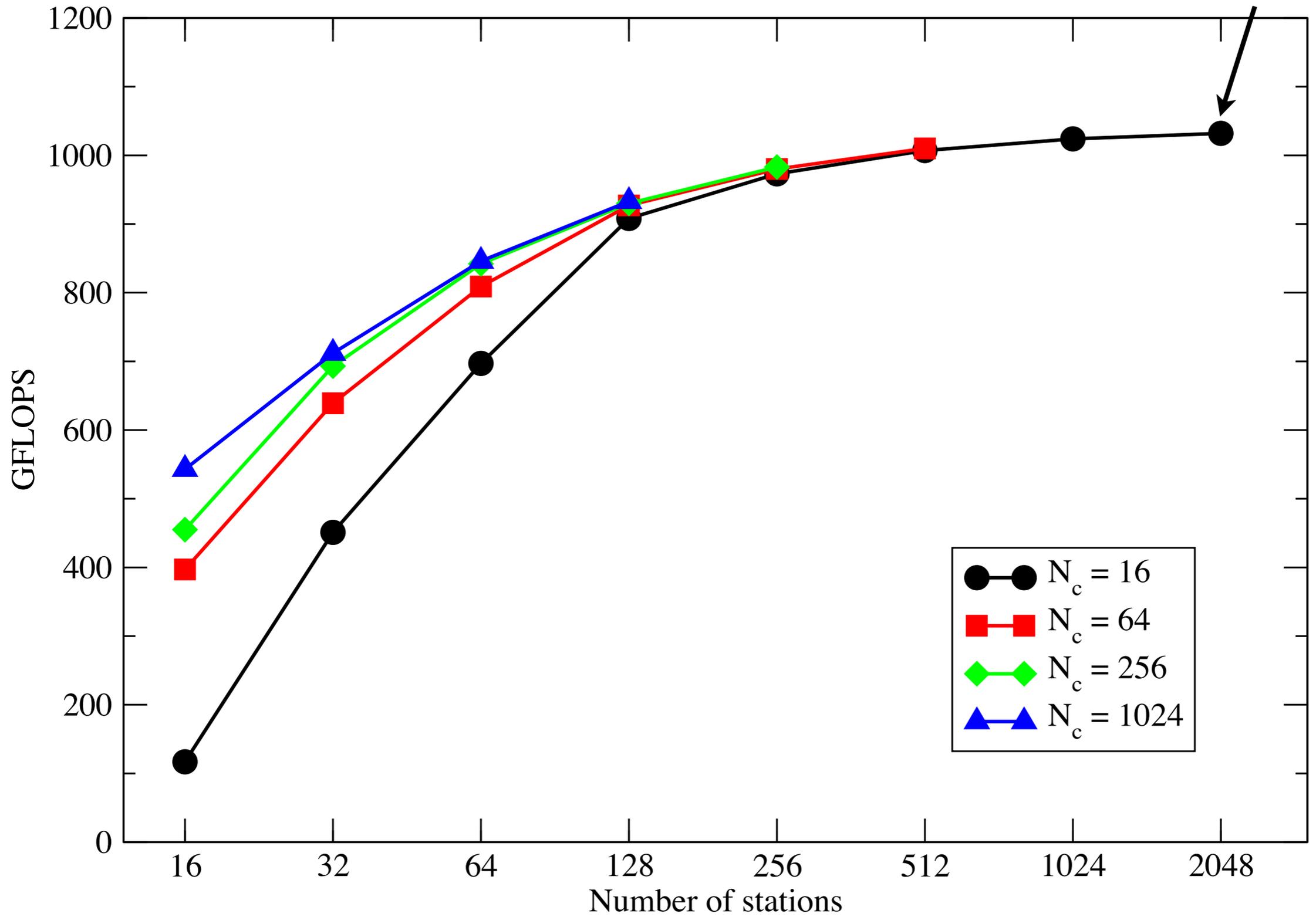


# Performance for $N_s=512$ , $N_c=12$ , $N_t=1000$





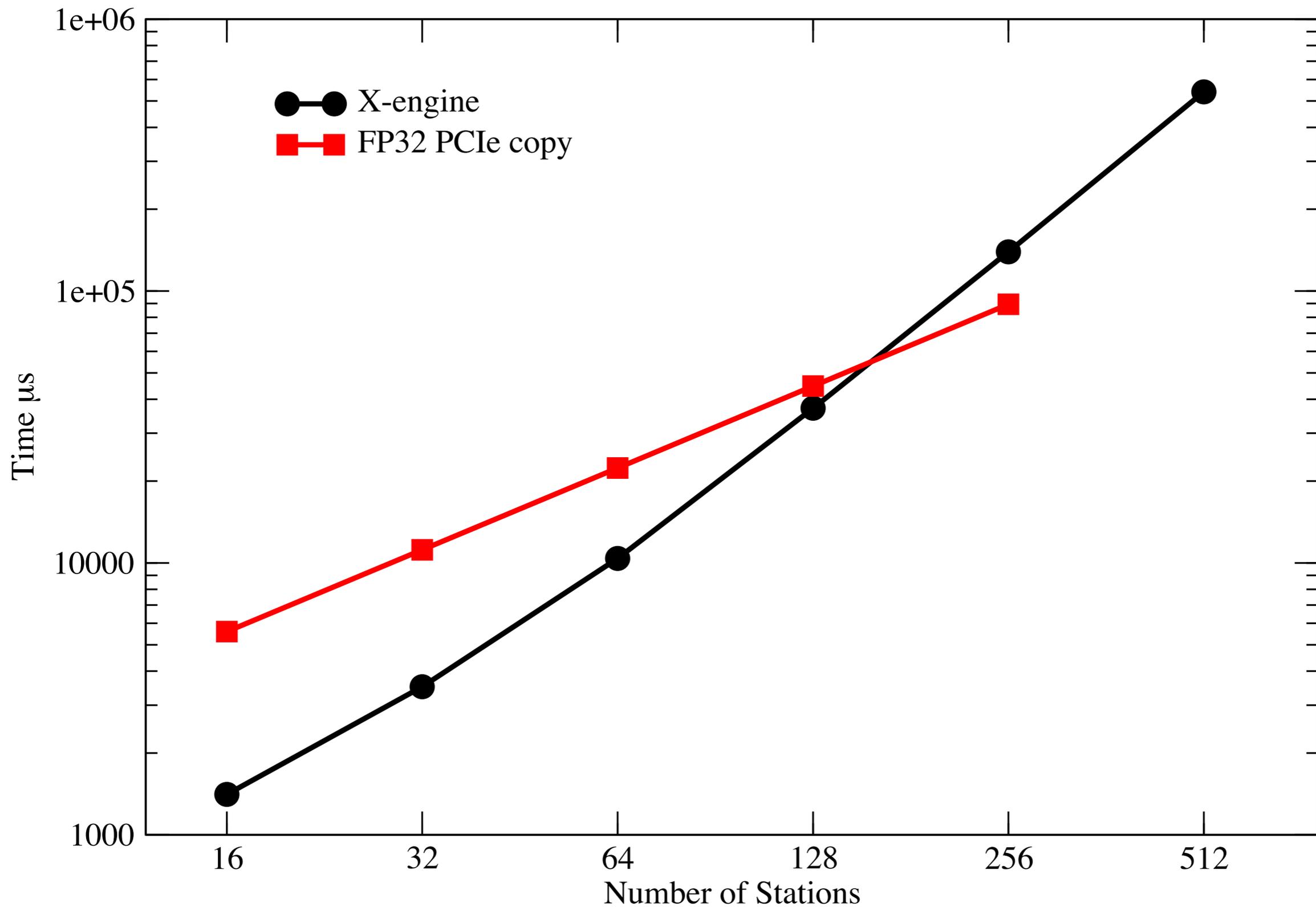
1042 GFLOPS = 78% (91%) peak  
67 GB/s = 38% peak



# Feeding the Beast

- 1 TFLOP sustained is all very well, but what if the bus can't keep up?
- Algorithm profile
  - Host  $\rightarrow$  Device  $O(N_c N_s B_c)$
  - Kernel  $O(N_c N_s^2 B_c)$
  - Device  $\rightarrow$  Host  $O(N_c N_s^2)$
- Kernel will dominate at large  $N_s$  but what about small  $N_s$ ?
- Previous work showed X-engine limited by bus
  - ~250 GFLOPS at  $N_s = 64$  (van Nieuwpoort *et al*)

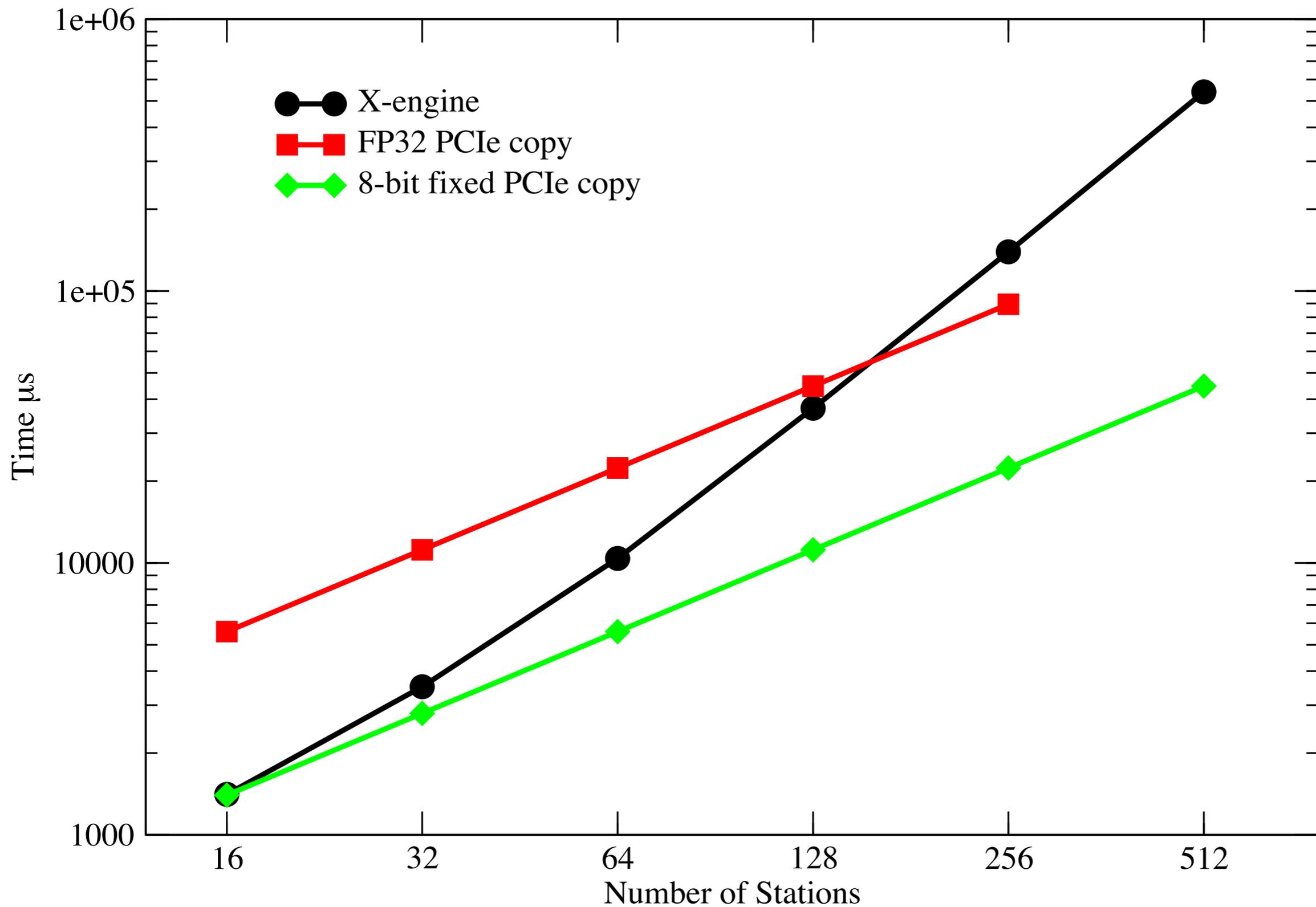
# Download time versus Kernel execution time



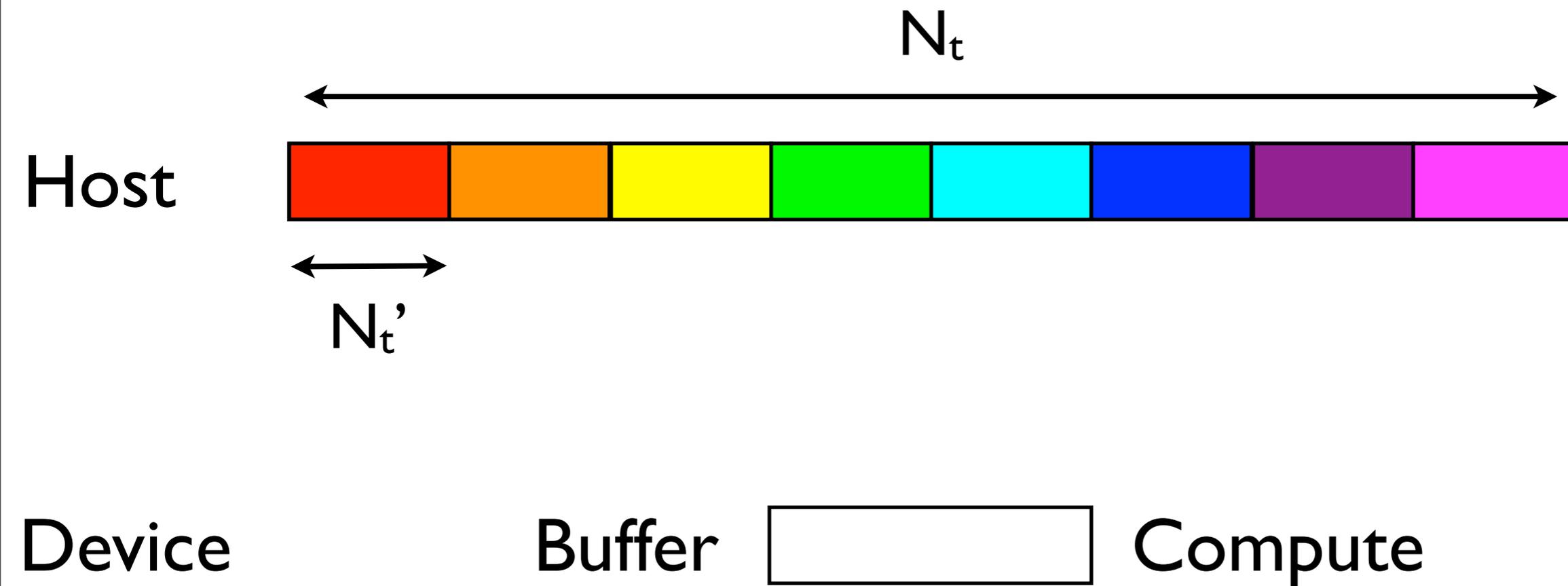
# The beast eats a lot, but each mouthful is small...

- PCIe bus is a severe constraint to performance
  - $N_s \geq 256$  kernel dominates
- Input signal precision typically 4-5 bits
- FP32 is a complete waste of bandwidth
- Accumulated correlation matrix must be high precision ( FPGAs typically use  $> 12$ -bit precision)
- Use 8-bit precision for input, keep matrix FP32
  - 8-bit natively supported through textures

# Download time versus Kernel execution time

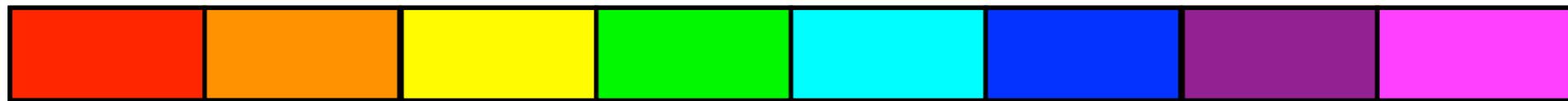


# Pipelined Algorithm



# Pipelined Algorithm

Host



Device

Buffer



Compute

# Pipelined Algorithm

Host



Device

Buffer



Compute

# Pipelined Algorithm

Host



Device

Buffer



Compute

# Pipelined Algorithm

Host



Device

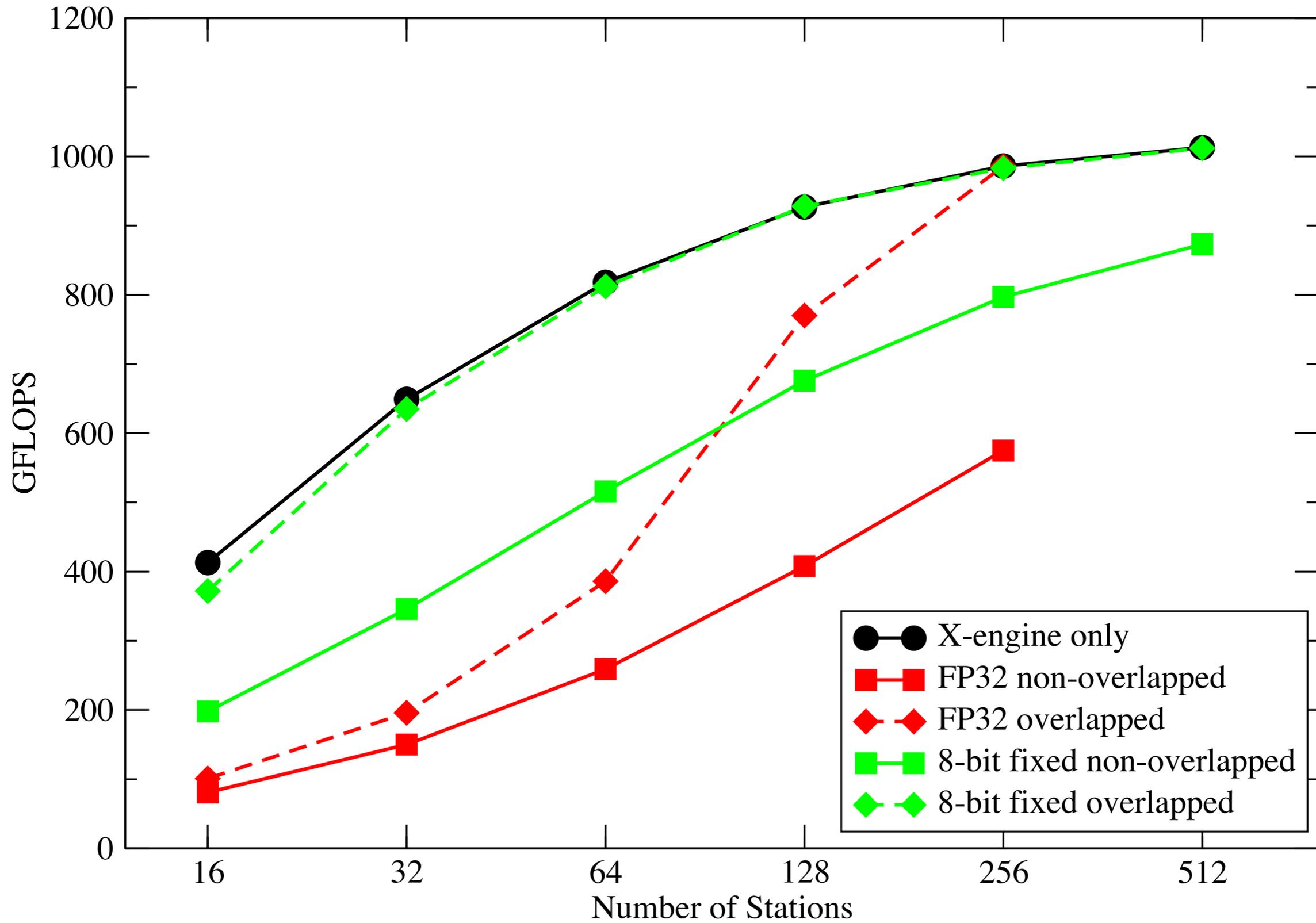
Buffer



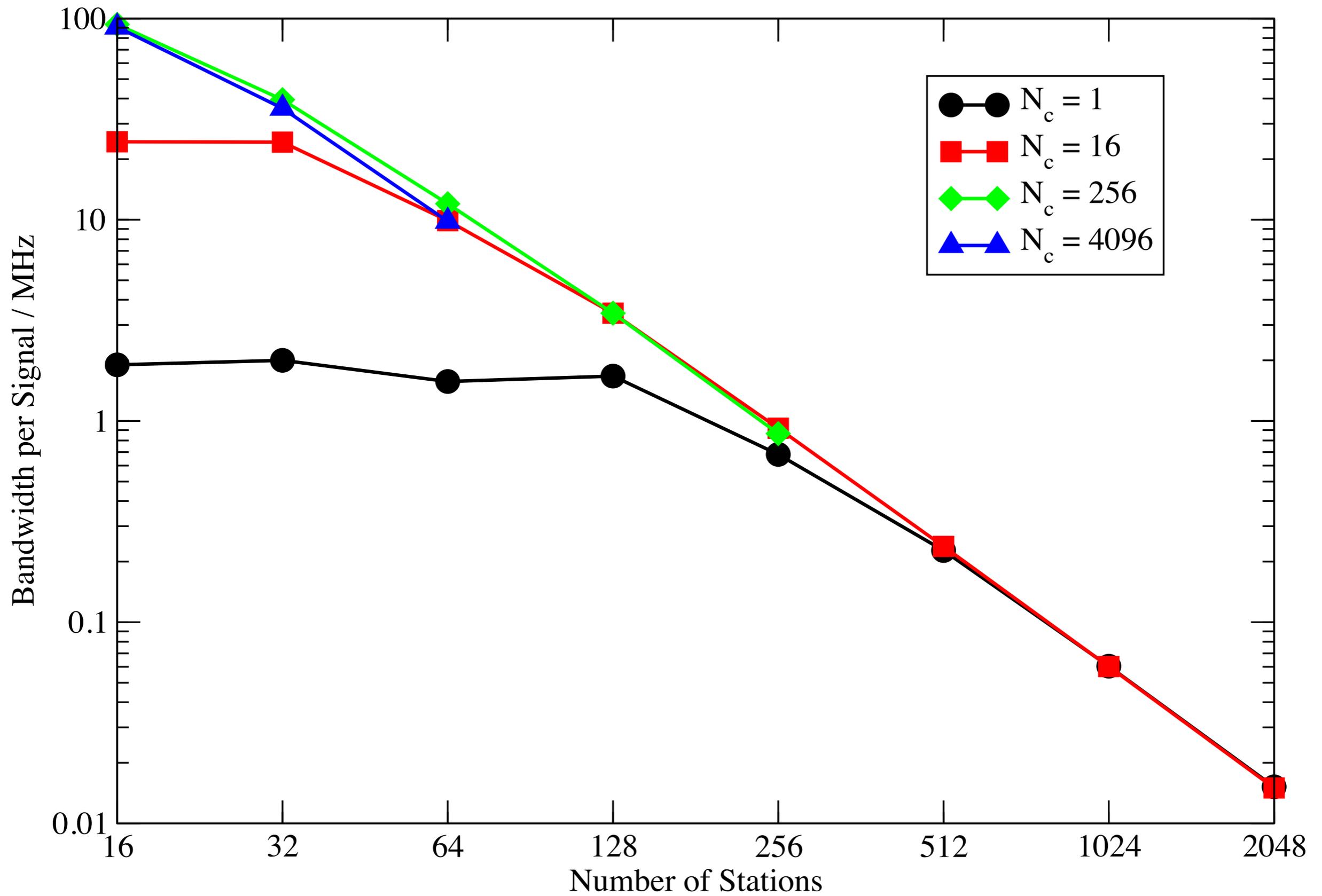
Compute

# Sustained X-engine performance

128 frequency channels, 1024 time samples



# Bandwidth achievable per signal





# Multi-GPU

- Trivial to parallelize across frequency
  - PCIe bus contention manageable
- 4 TFLOPs sustained with 4 GPUs
  - 1300 Watts total
    - 3 GFLOPS/Watt (cf 1684 GFLOPS/Watt #1 Green 500)
  - \$6K for workstation => \$1.50 per GFLOP

# X-engine Performance Across Platforms

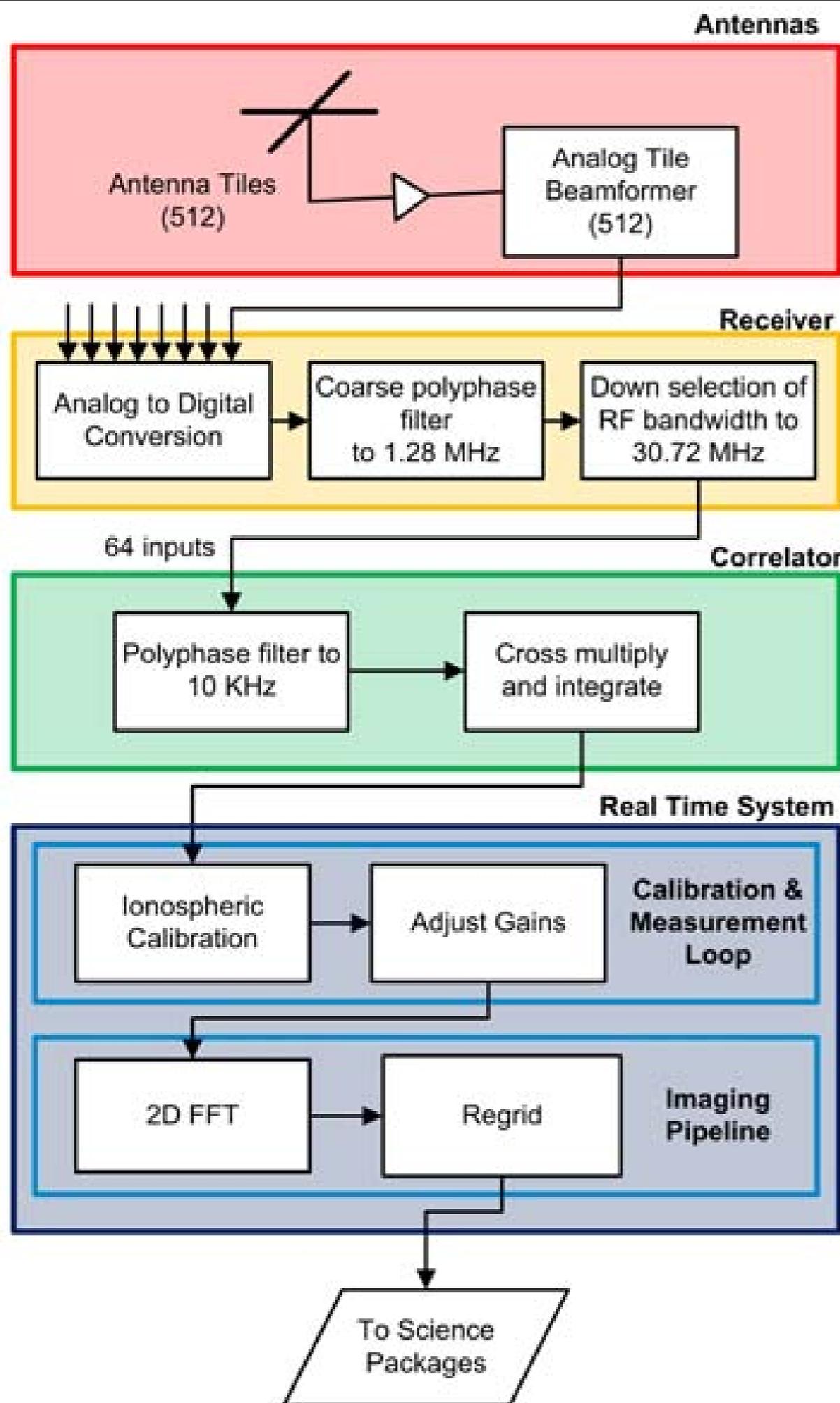
Architecture	GFLOPS	GOPS/Watt per chip (total)
Intel Core i7 (quad)*	48.0	0.4
IBM BG/P*	13.1	0.5
IBM Cell*	187	2.7
Nvidia C1060*	243	1
Nvidia GTX 480†	1042	~4 (3)

\* van Nieuwpoort and Romein, †this work, #de Souza et al and Lonsdale et al, †Manley

# X-engine Performance Across Platforms

Architecture	GFLOPS	GOPS/Watt per chip (total)
Intel Core i7 (quad)*	48.0	0.4
IBM BG/P*	13.1	0.5
IBM Cell*	187	2.7
Nvidia C1060*	243	1
Nvidia GTX 480†	1042	~4 (3)
MWA Correlator # (Virtex 4 SX35)		(~10)
Roach II+ (Virtex 6)		(~58)

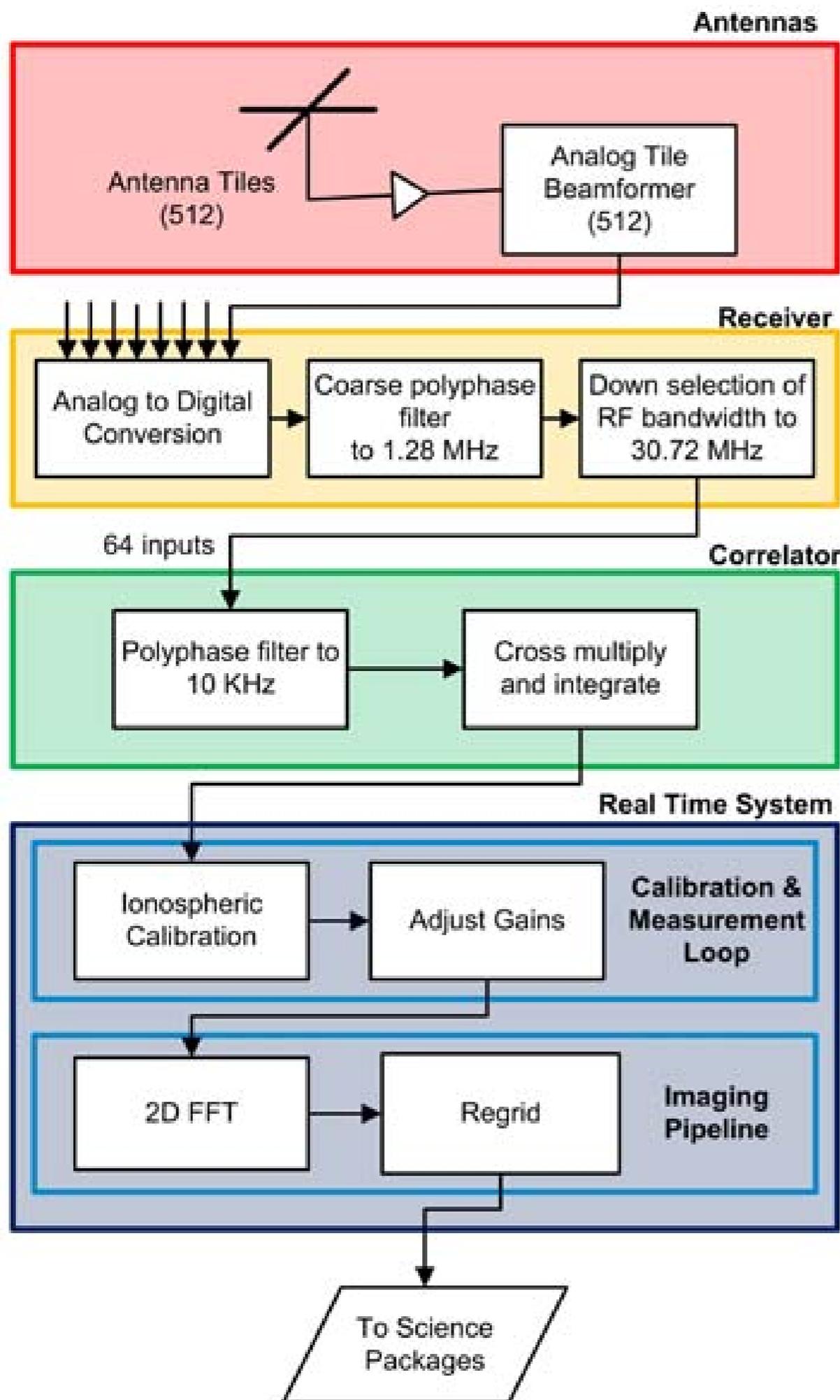
\* van Nieuwpoort and Romein, †this work, #de Souza et al and Lonsdale et al, +Manley



FPGAs  
Virtex 4 SX35 16 kW

64 x GPU  
cluster 30 kW

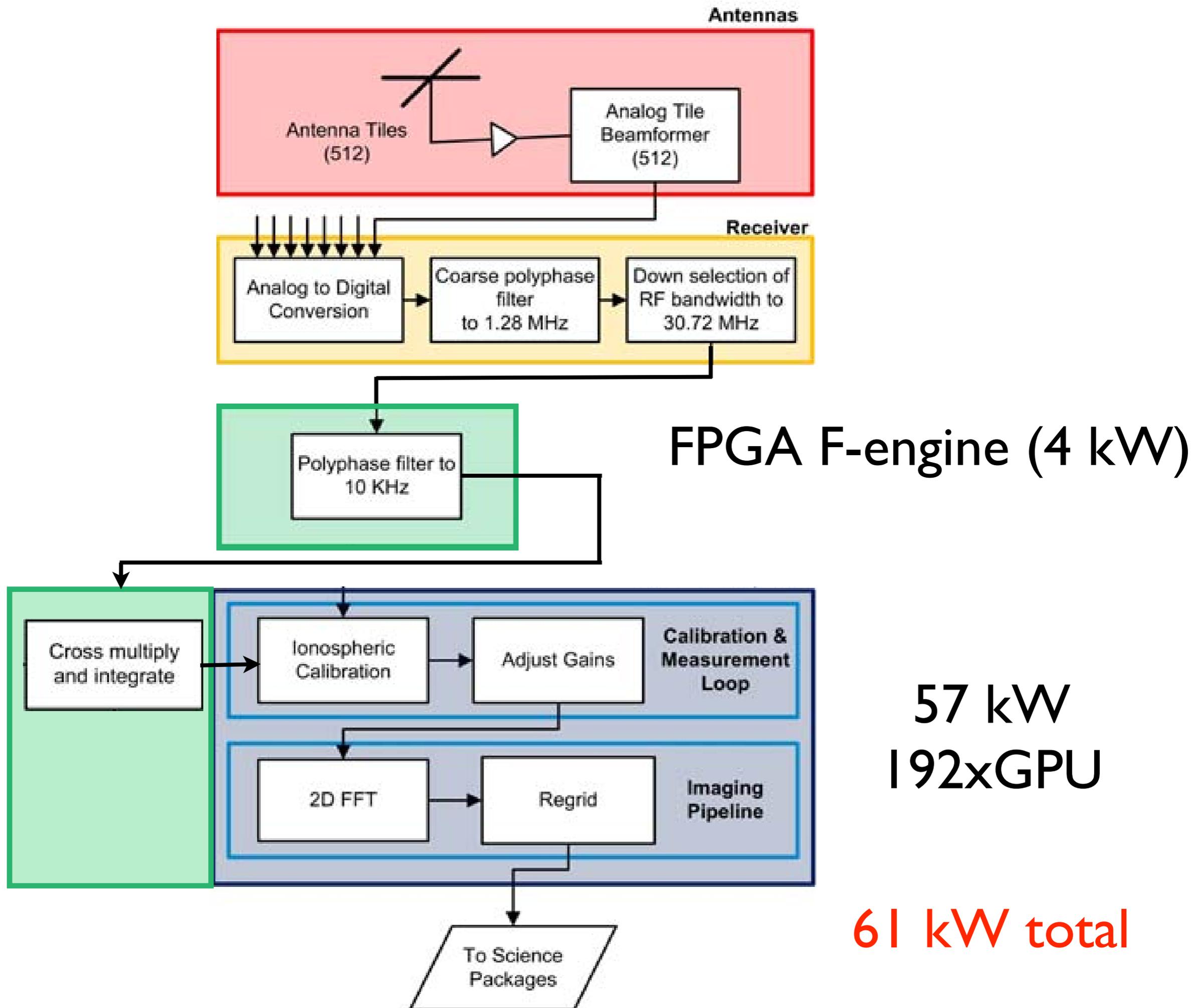
46 kW total



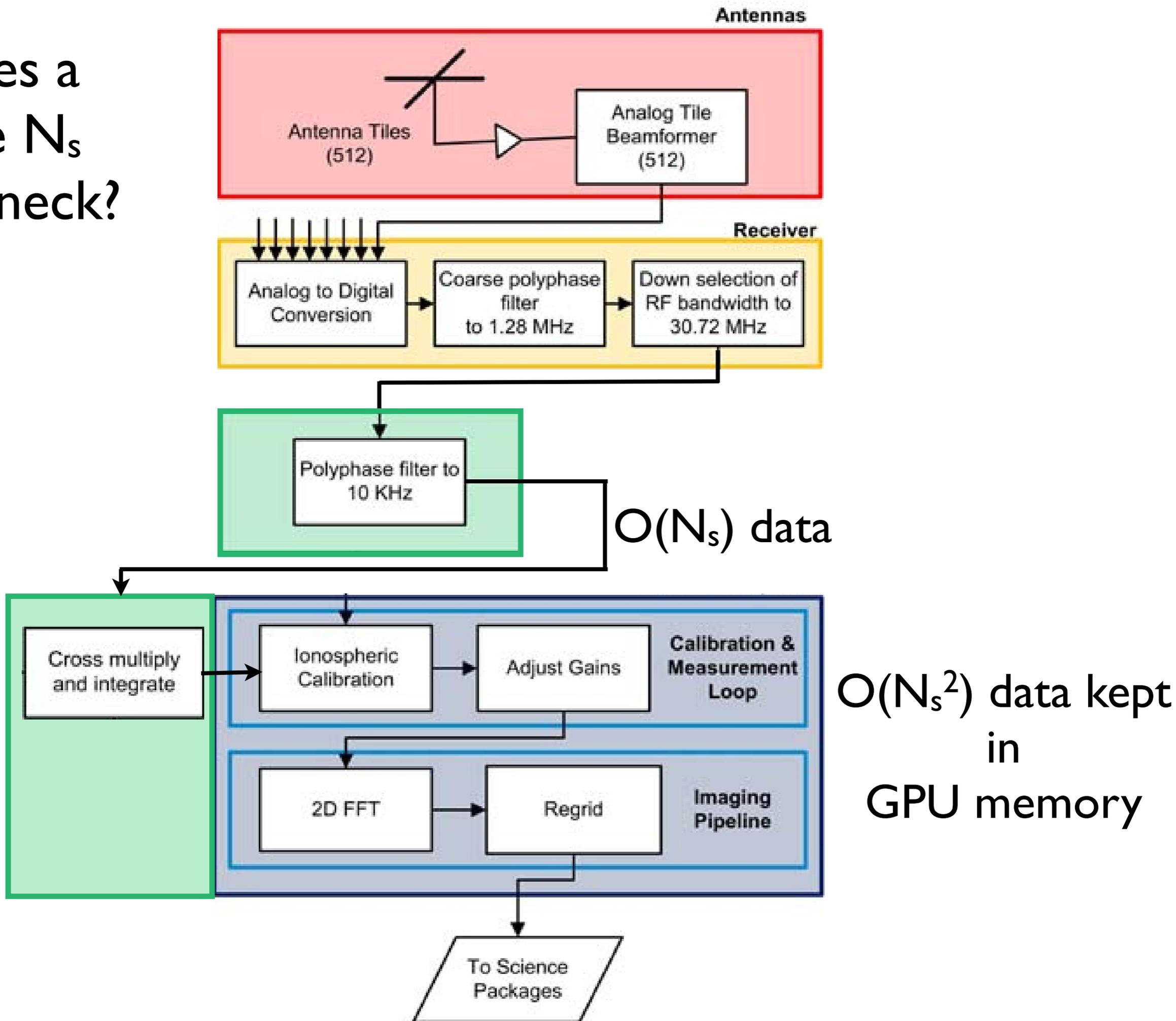
FPGA F-engine 4 kW  
 128x GPU X-engine 42 kW

64x GPU 30 kW

76 kW total

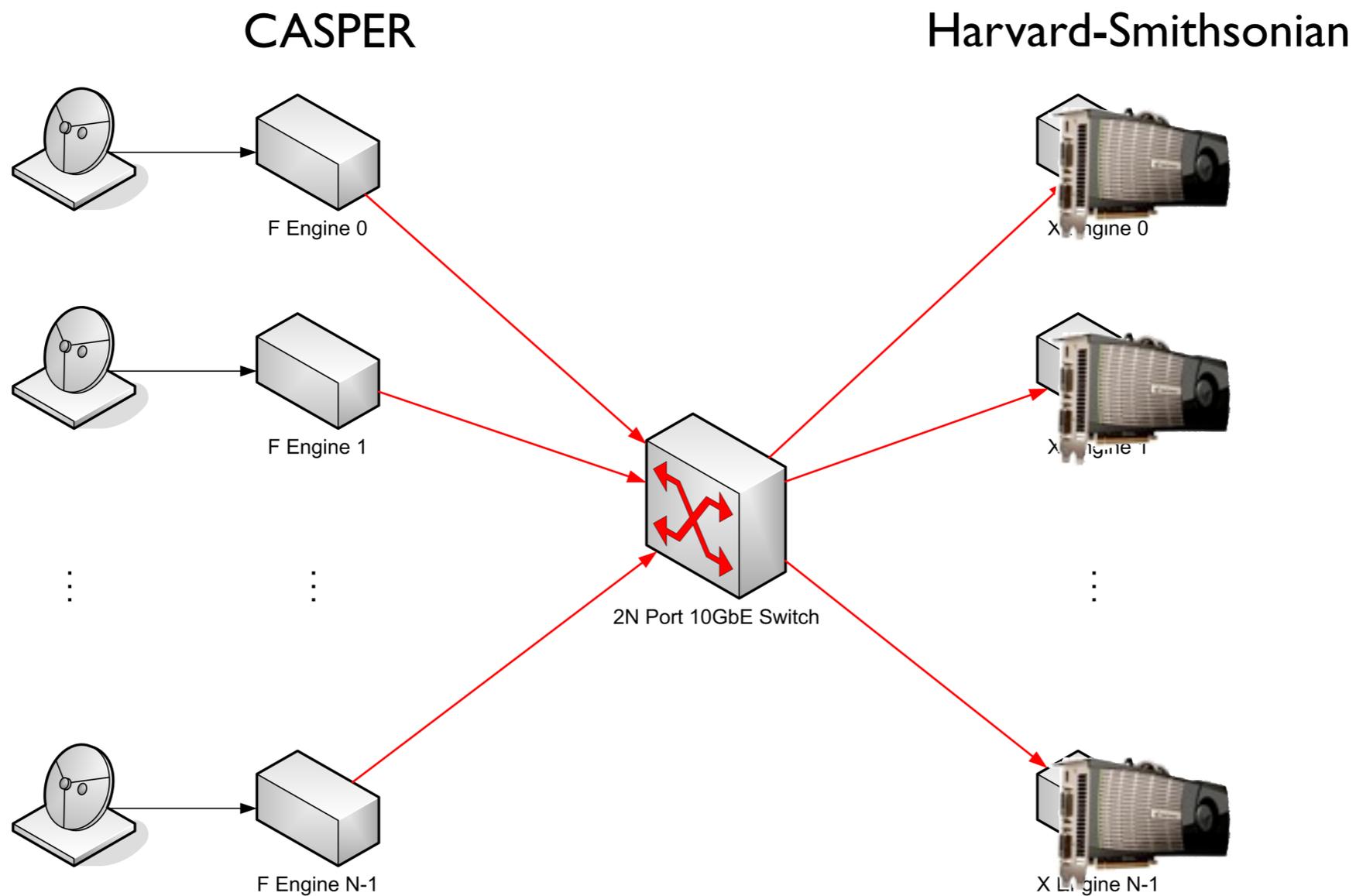


Solves a large  $N_s$  bottleneck?

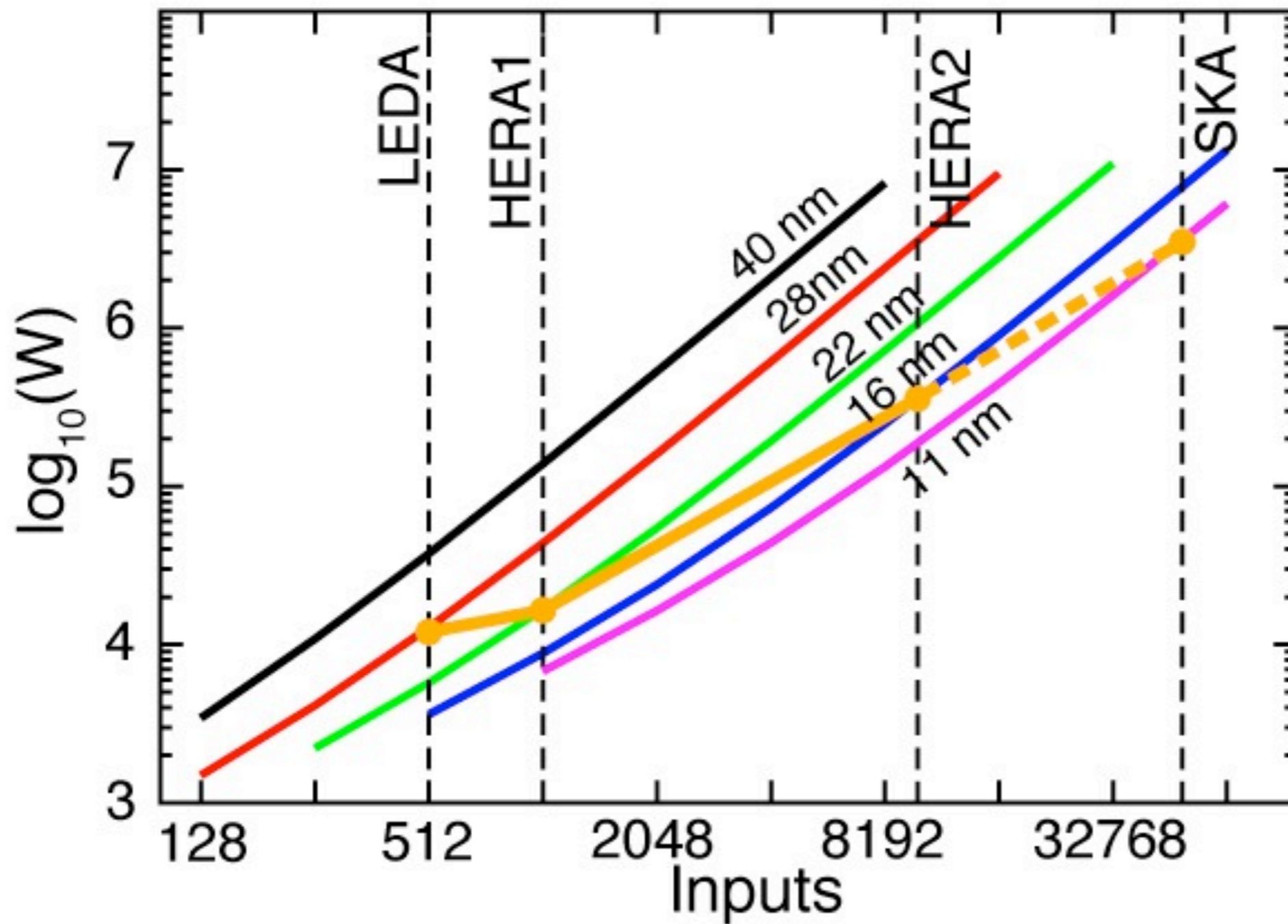


# Hybrid Correlator

- Proposed correlator for LEDA



# Scaling to the Future



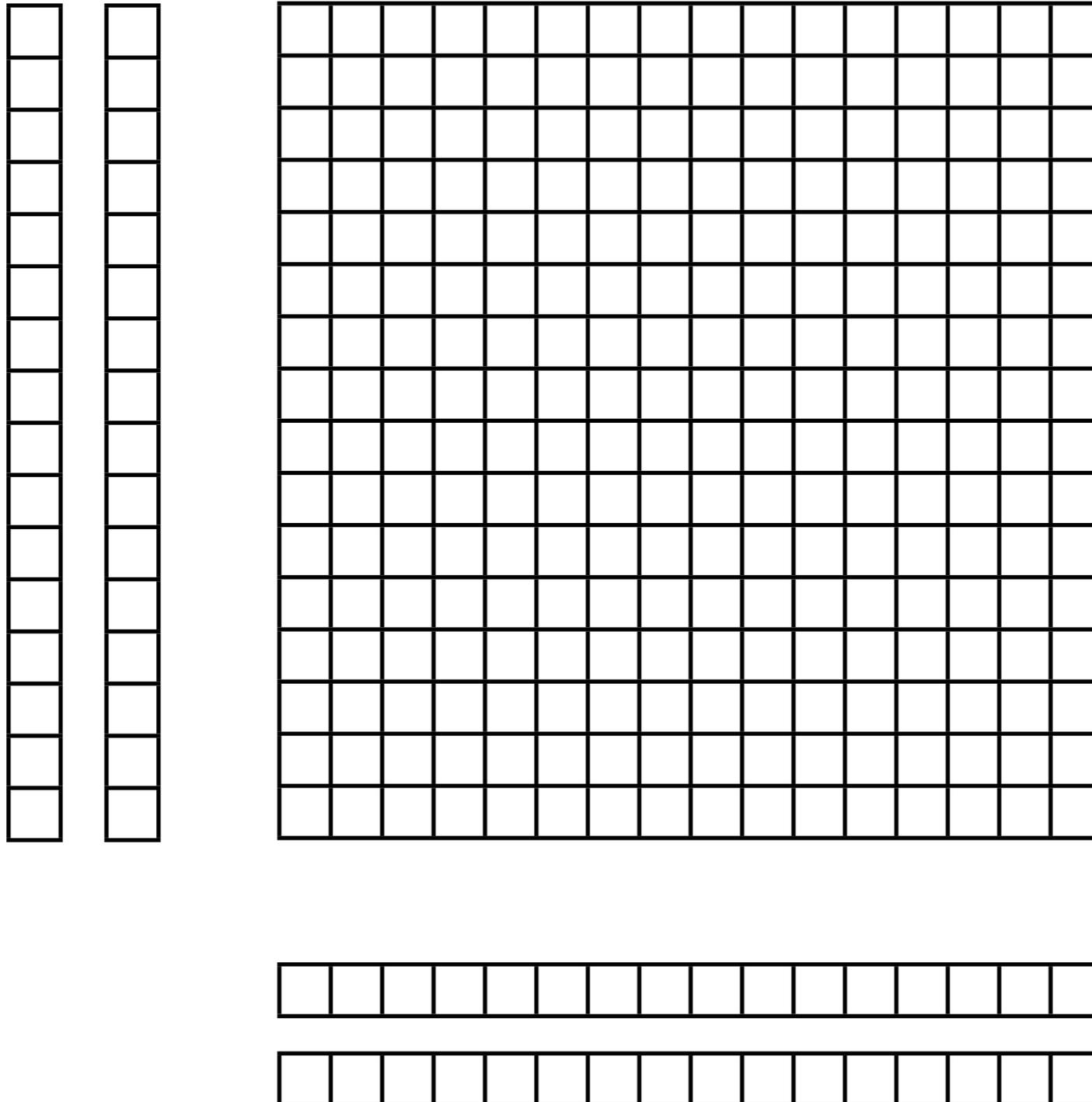
# Summary and Conclusion

- X-engine is a perfect match to the GPU
  - 78% peak performance
- Low cost and development time
  - Easy to keep with bleeding edge
- Not (yet) power competitive with FPGAs
- Future: hybrid correlators?
- Combine X-engine with Calibration and Imaging

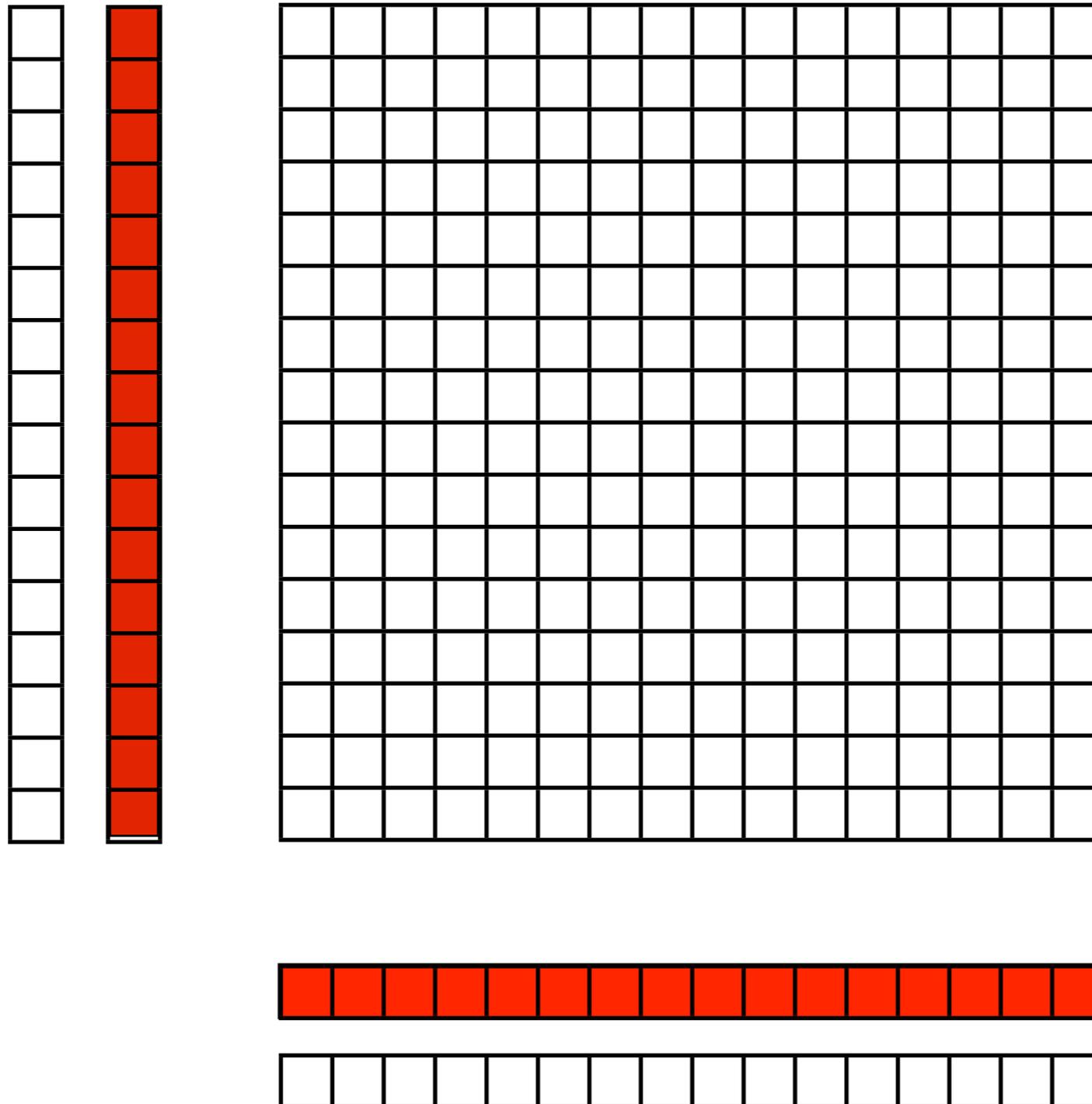
# Optimization lessons

- Shared memory AND register tiling critical
- Minimize integer arithmetic
  - Precalculate offsets
  - Texture lookup
- Thread synchronization is costly
- The compiler doesn't always know what it's doing

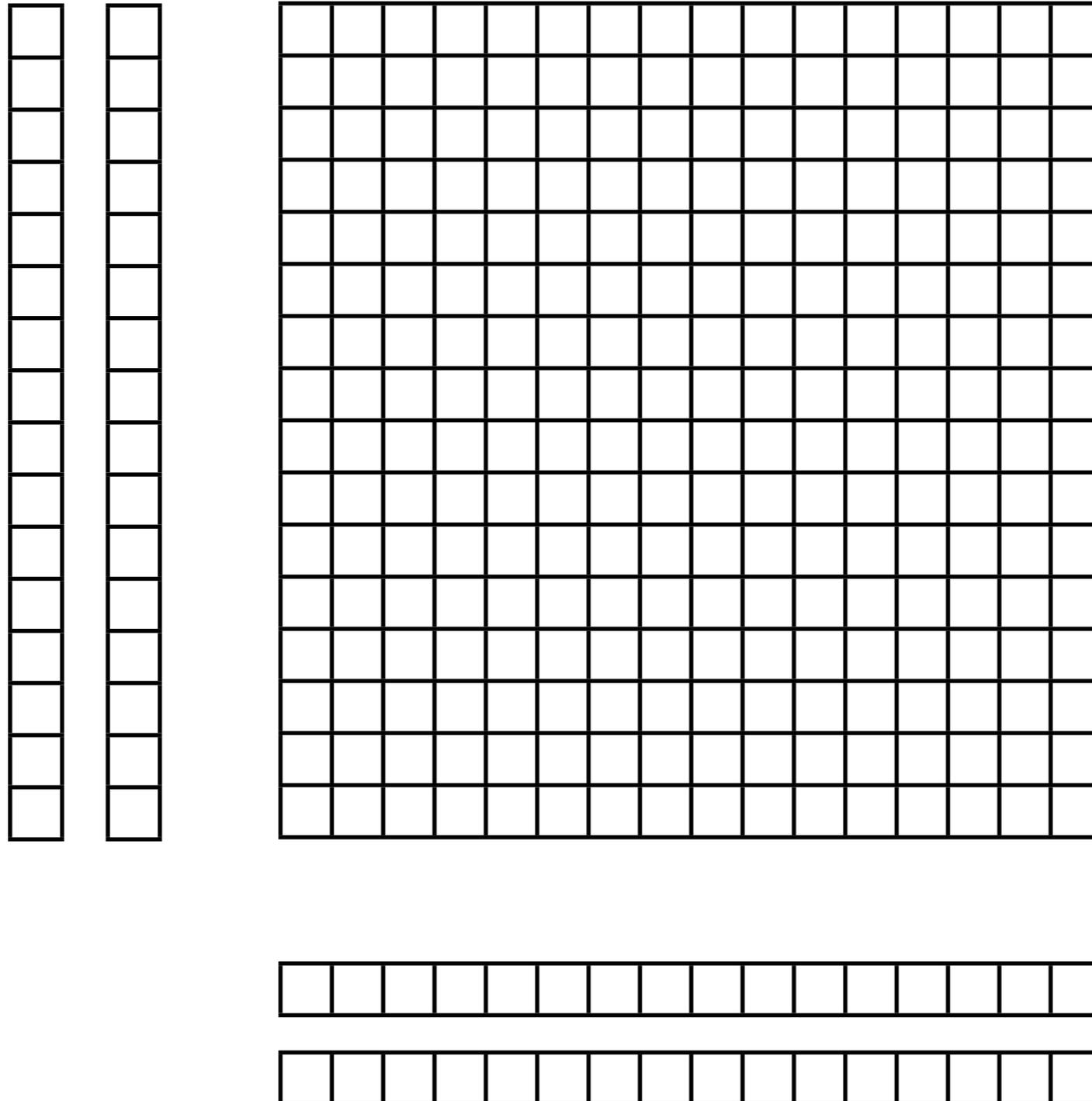
# Double buffer shared memory storage to reduce synchronization



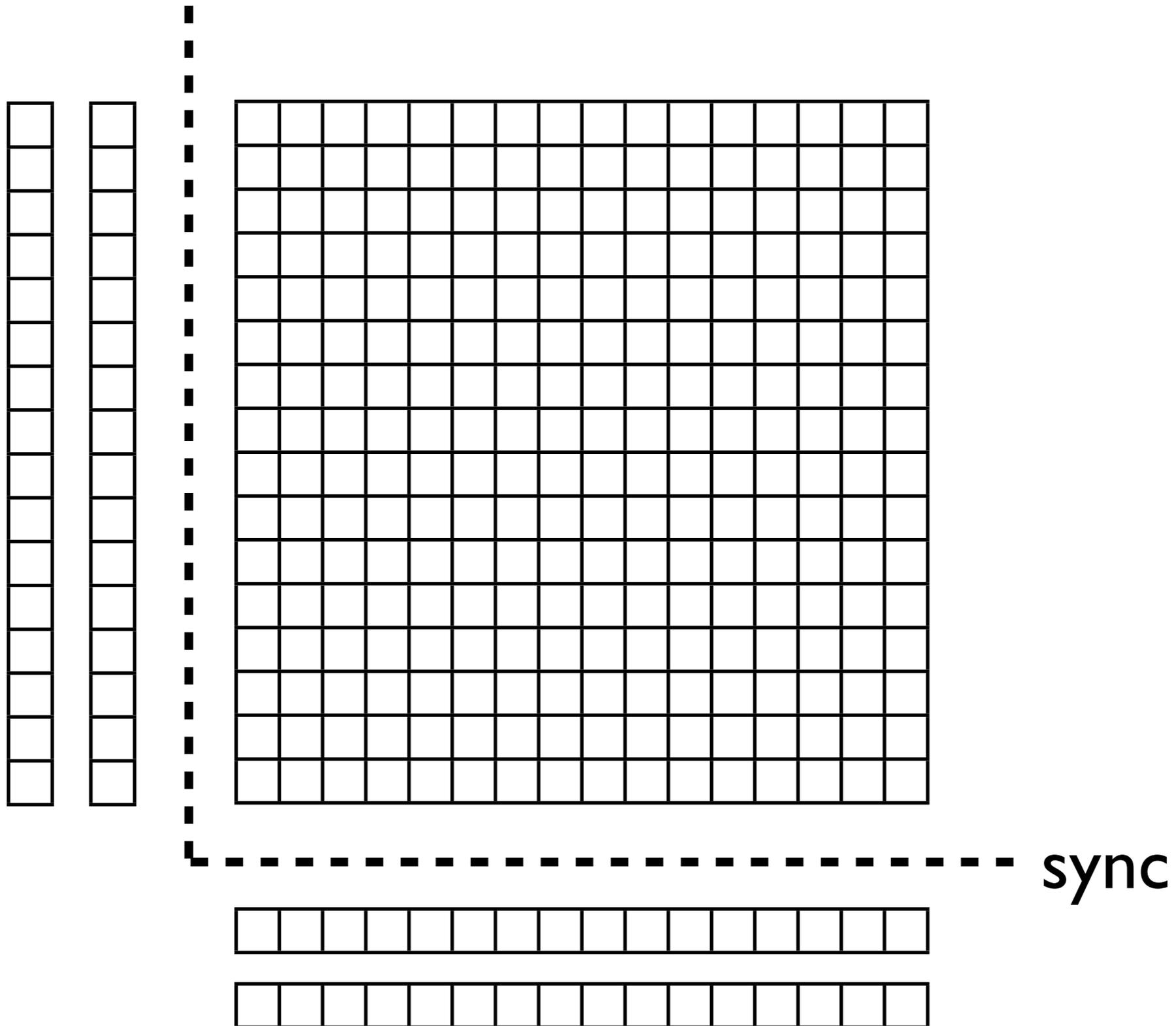
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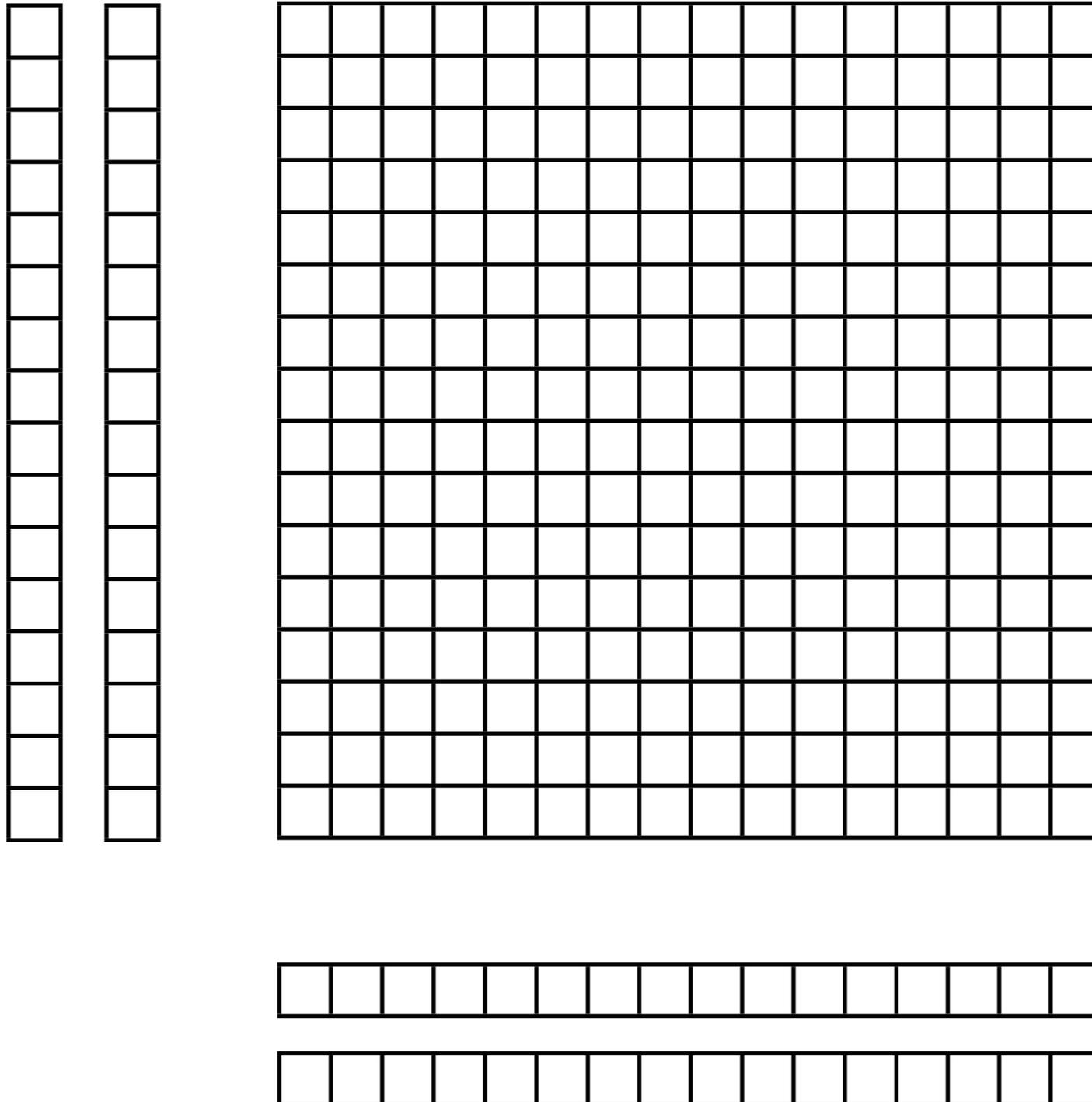
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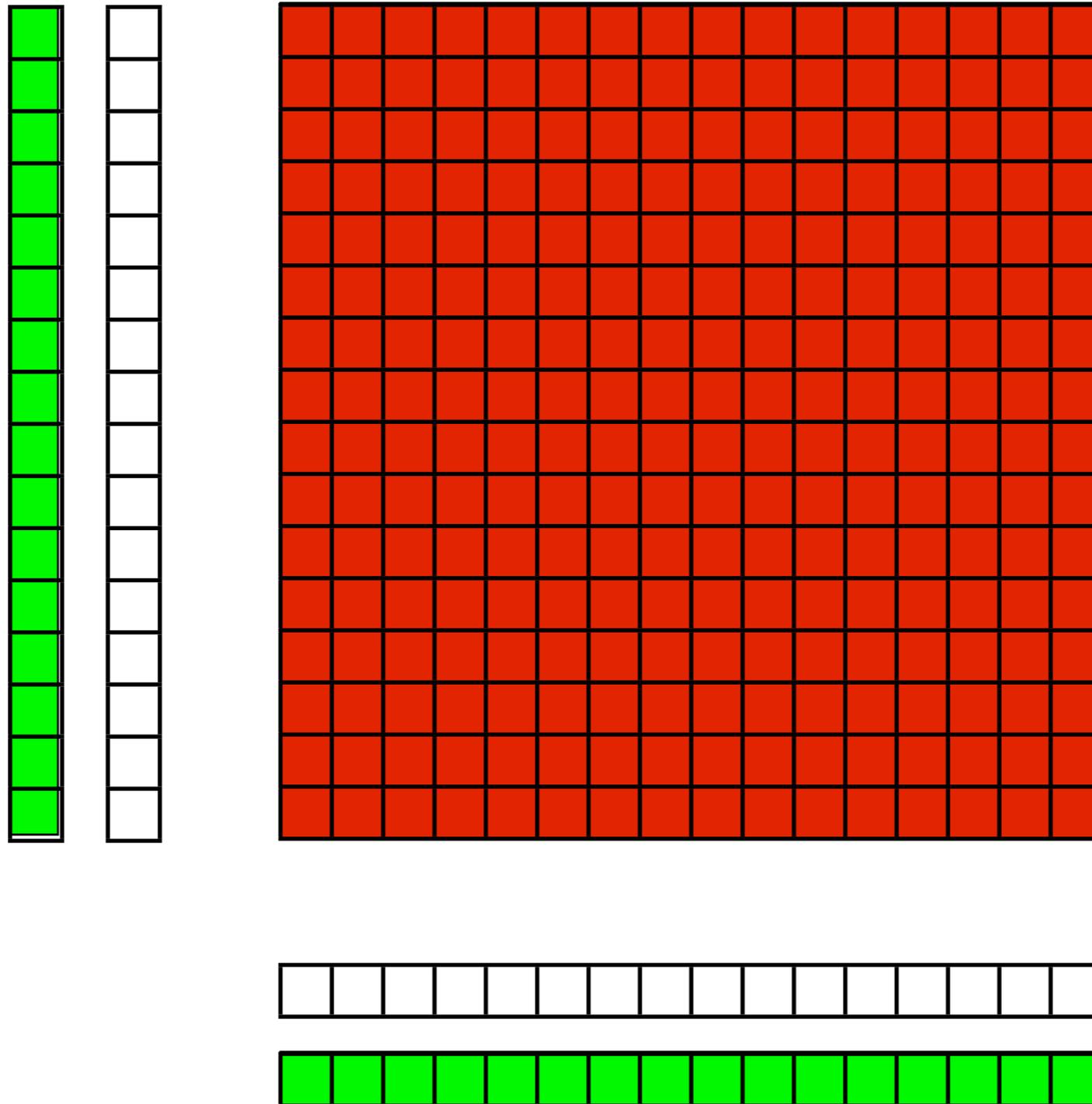
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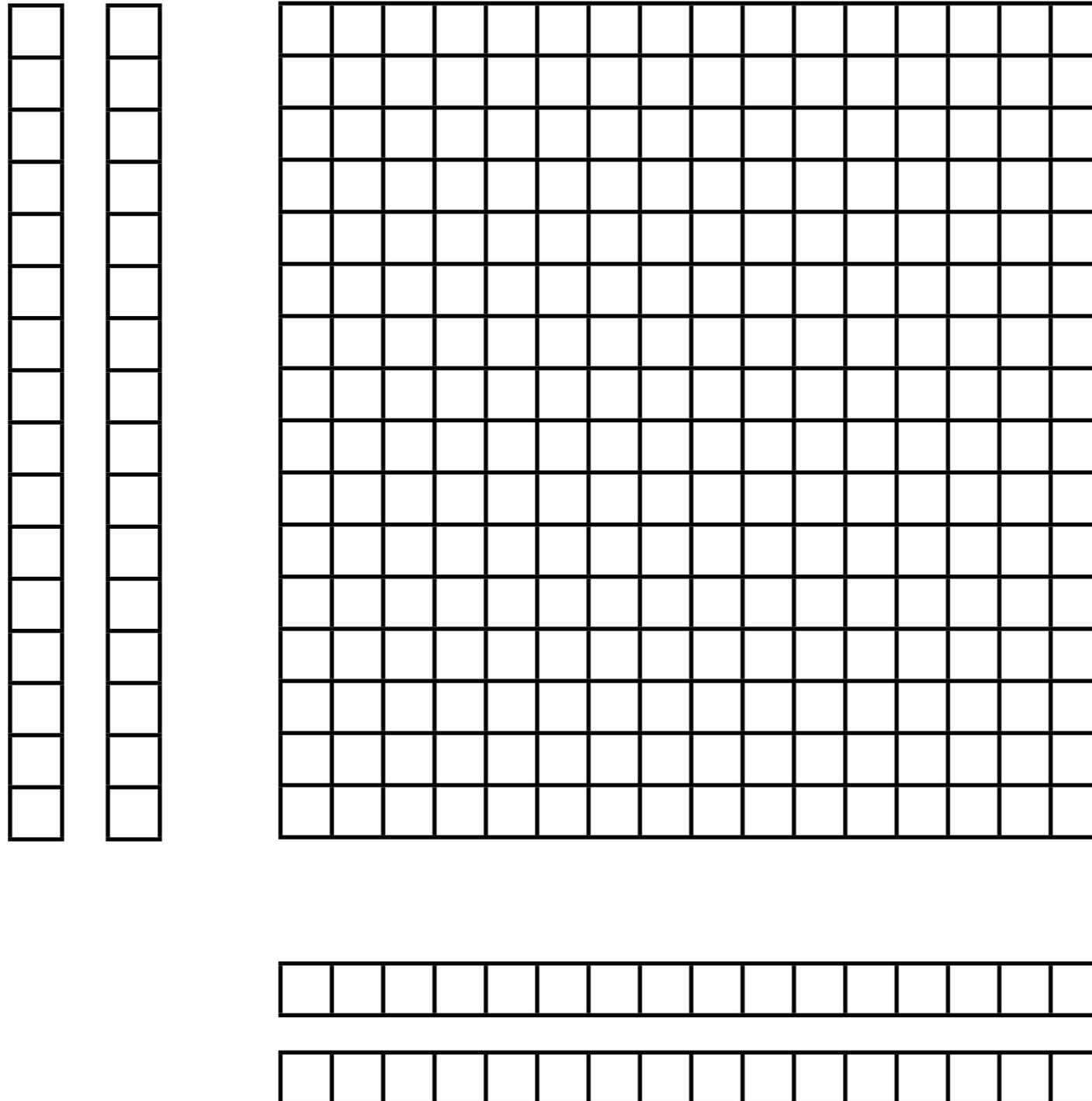
# Double buffer shared memory storage to reduce synchronization



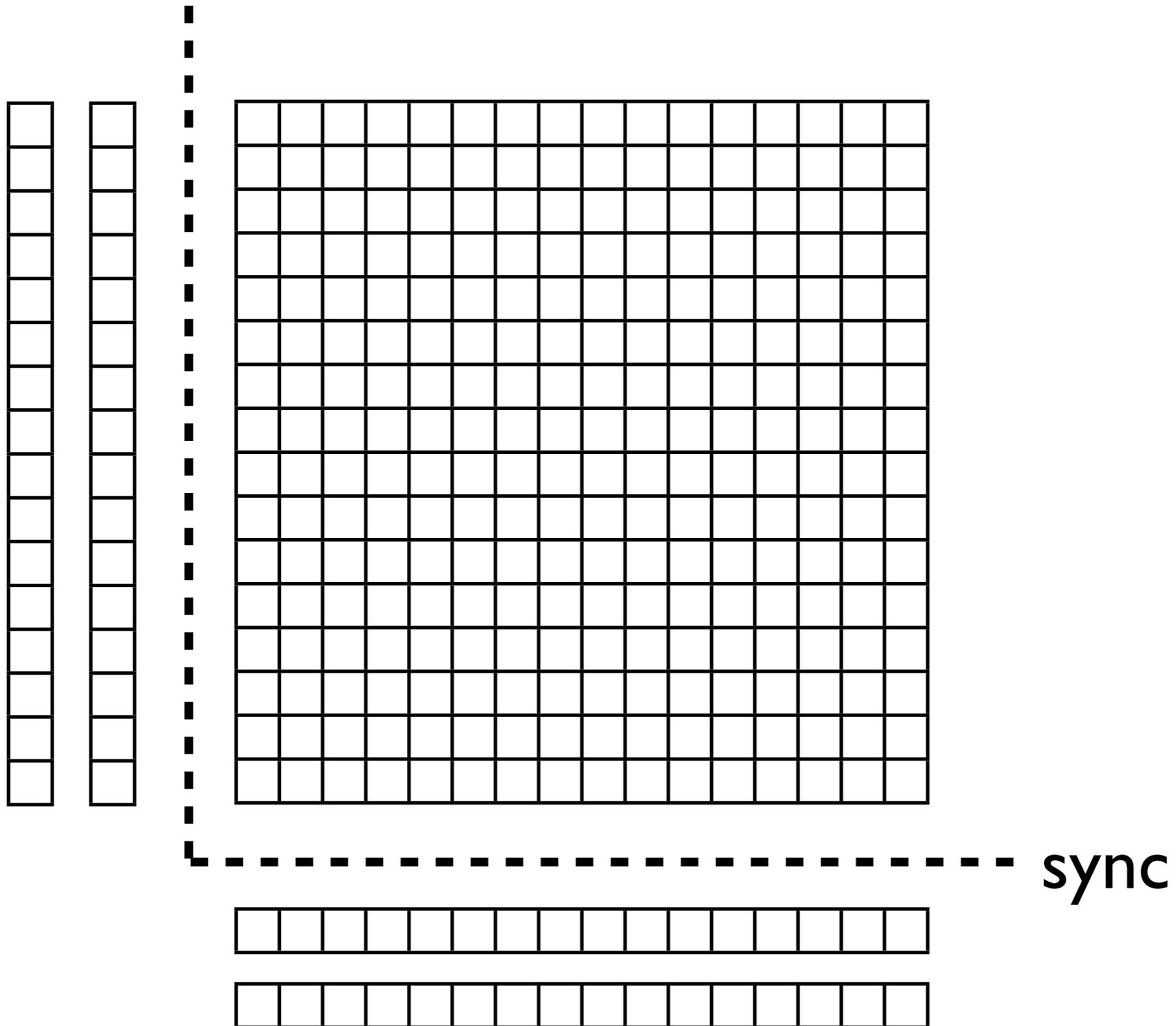
# Double buffer shared memory storage to reduce synchronization



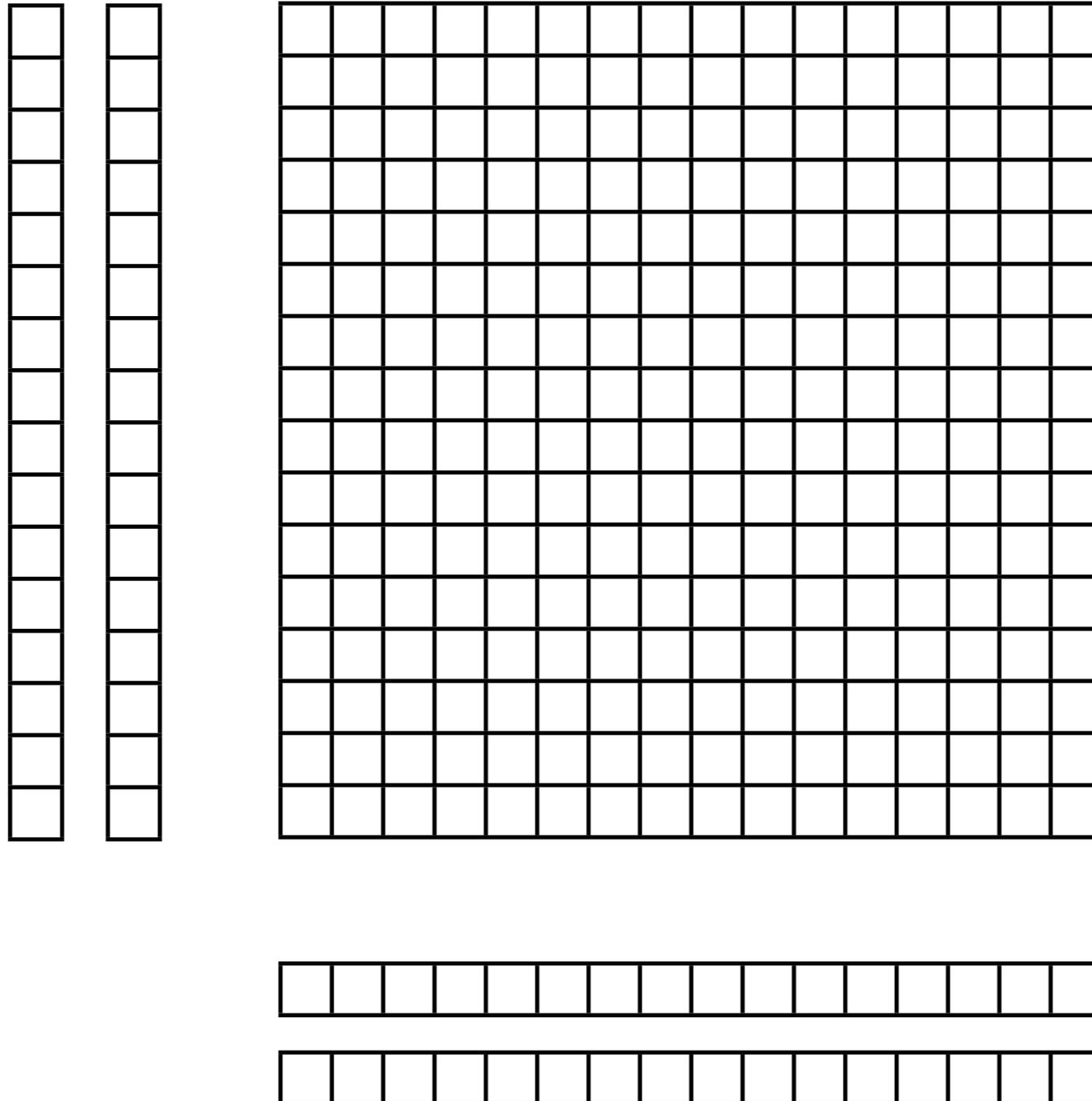
# Double buffer shared memory storage to reduce synchronization



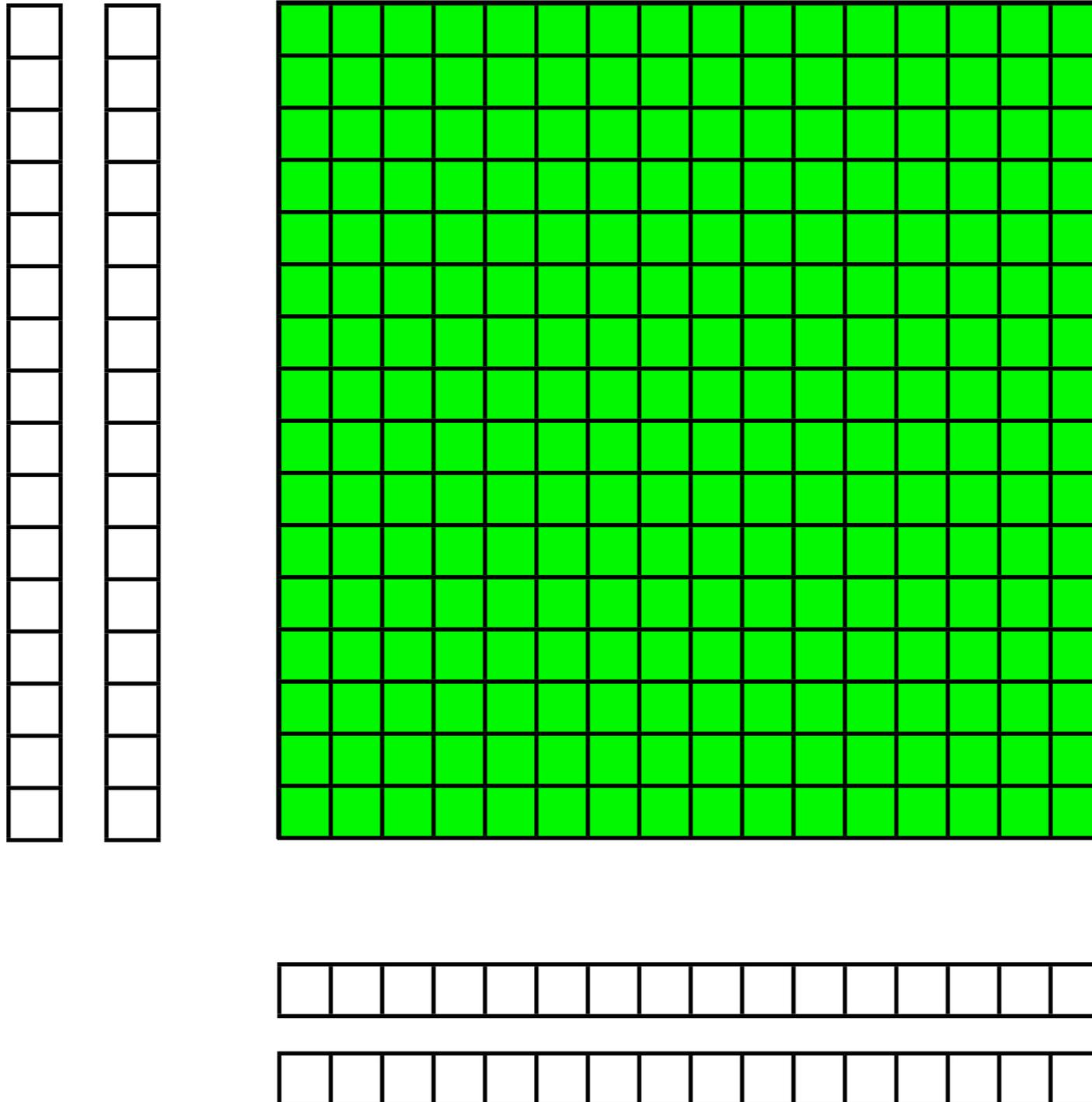
# Double buffer shared memory storage to reduce synchronization



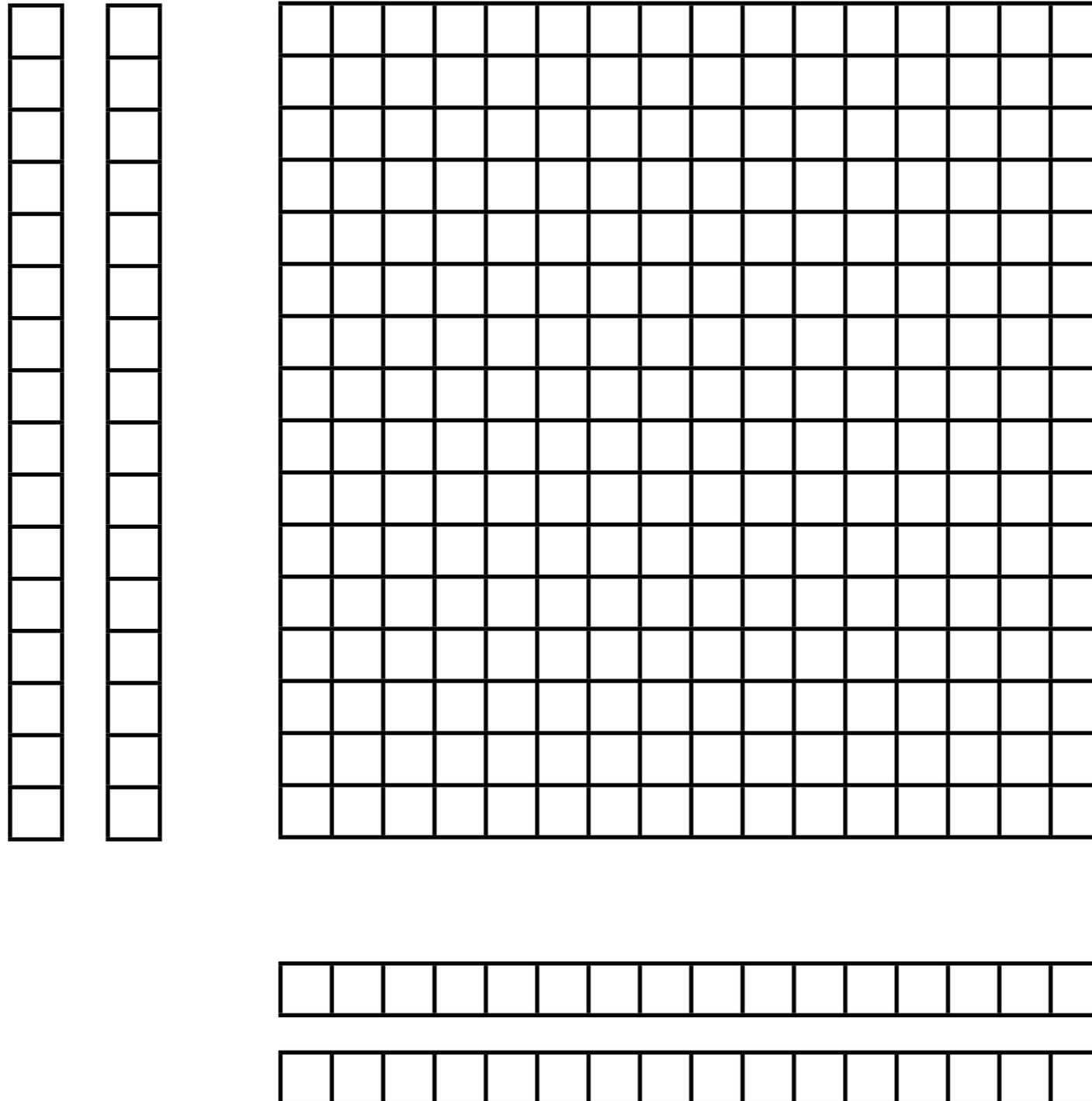
# Double buffer shared memory storage to reduce synchronization



# Double buffer shared memory storage to reduce synchronization



# Double buffer shared memory storage to reduce synchronization



# Don't trust compilers

- Compare these “identical” code fragments

```
a += b*c + d*c + e*f + g*h;
```

```
a += b*c;
```

770 GFLOPS

```
a += d*c;
```

```
a += e*f;
```

1020 GFLOPS

```
a += g*h;
```